

Derivation Tree (6A)

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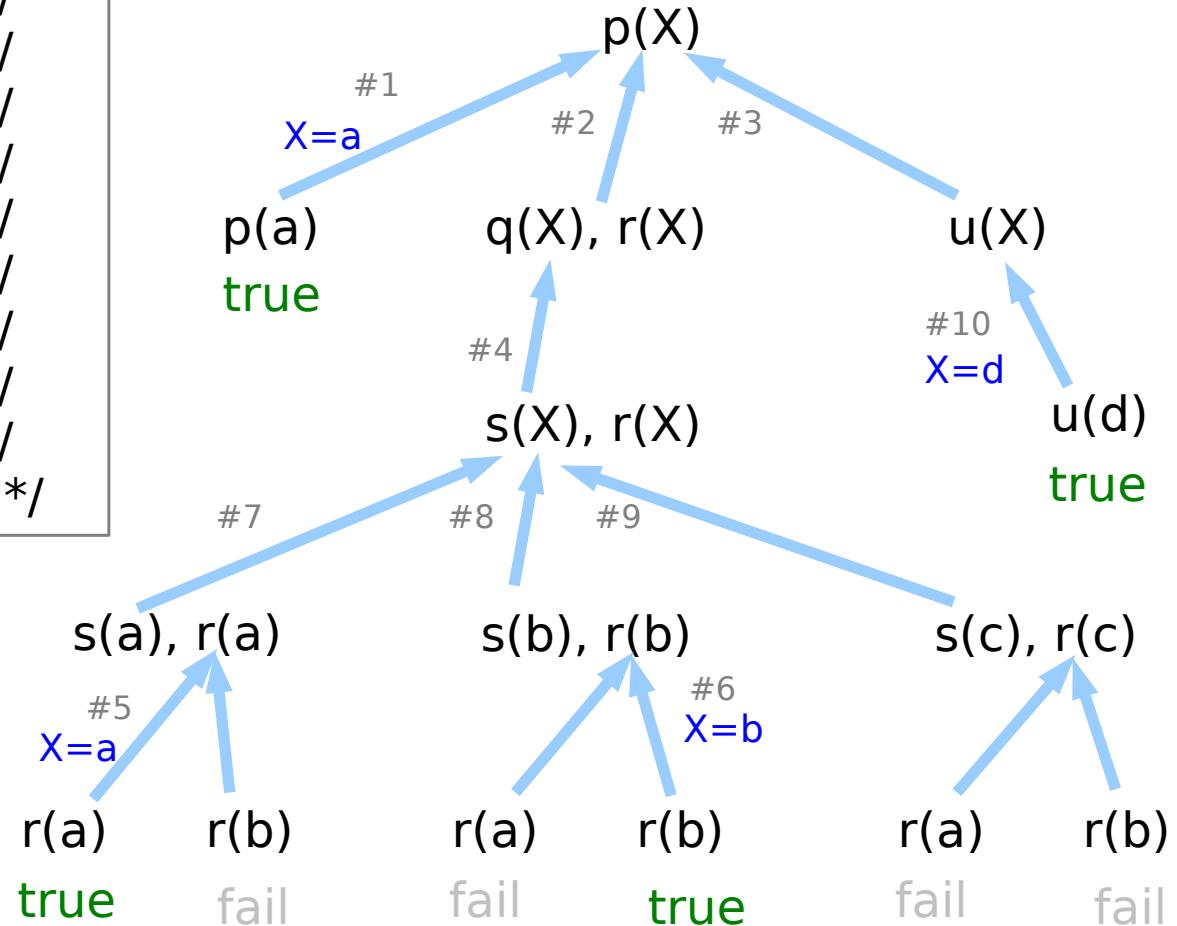
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Derivation Tree Examples

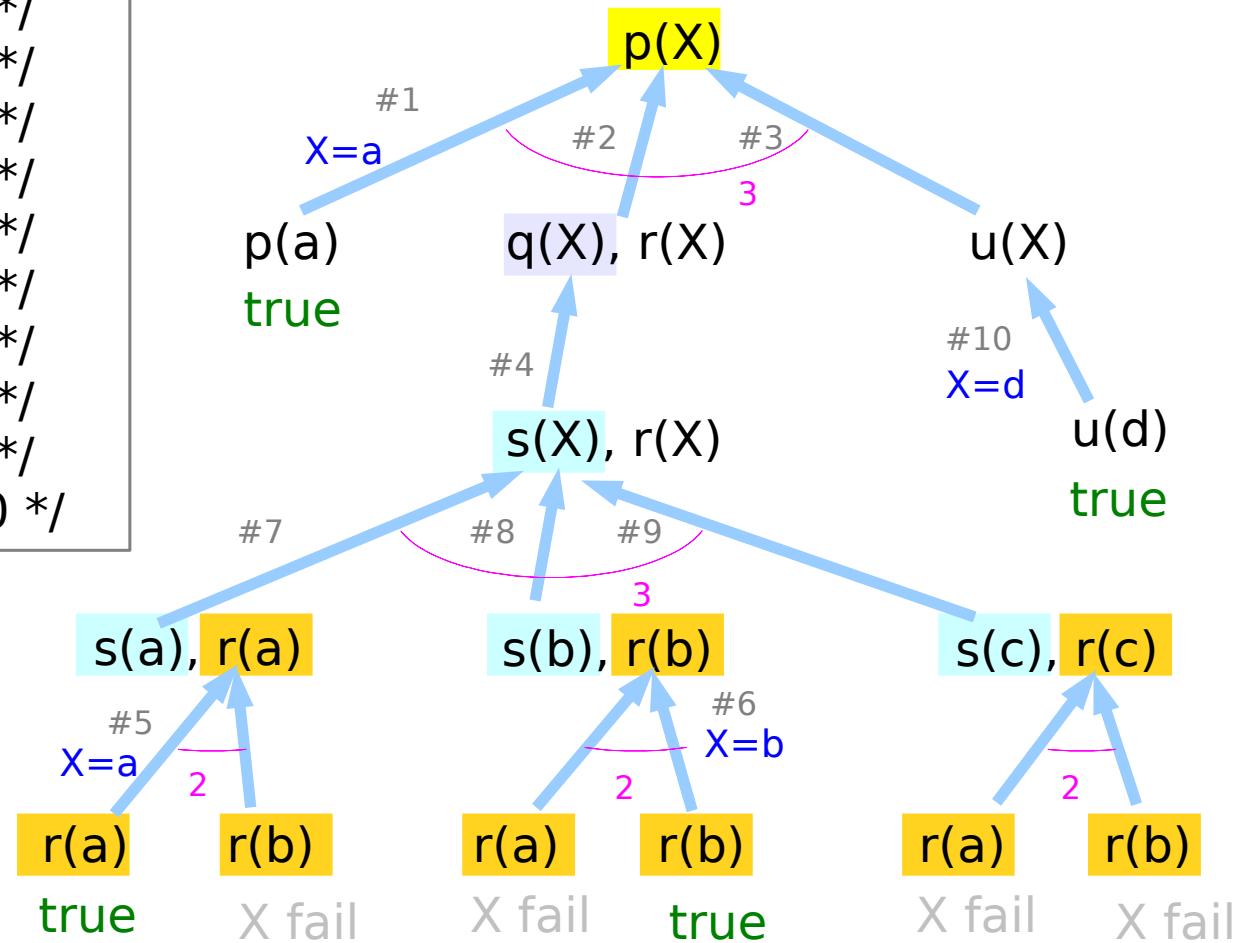
```
p(a).          /* #1 */  
p(X) :- q(X), r(X). /* #2 */  
p(X) :- u(X).    /* #3 */  
q(X) :- s(X).    /* #4 */  
r(a).          /* #5 */  
r(b).          /* #6 */  
s(a).          /* #7 */  
s(b).          /* #8 */  
s(c).          /* #9 */  
u(d).          /* #10 */
```



Derivation Tree Examples

```

3 { p(a). /* #1 */
    p(X) :- q(X), r(X). /* #2 */
    p(X) :- u(X). /* #3 */
    q(X) :- s(X). /* #4 */
1 {   r(a). /* #5 */
    r(b). /* #6 */
2 {   s(a). /* #7 */
    s(b). /* #8 */
    s(c). /* #9 */
3 {   u(d). /* #10 */
    
```



Replacing a selected subgoal

each node
the current goal
a sequence of subgoals

edges
the choices available for
replacing a selected subgoal

when g_1 unifies with h

g_1, g_2, g_3, \dots

$h :- b_1, b_2, \dots, b_n$

$b_1, b_2, \dots, b_n, g_2, g_3, \dots$

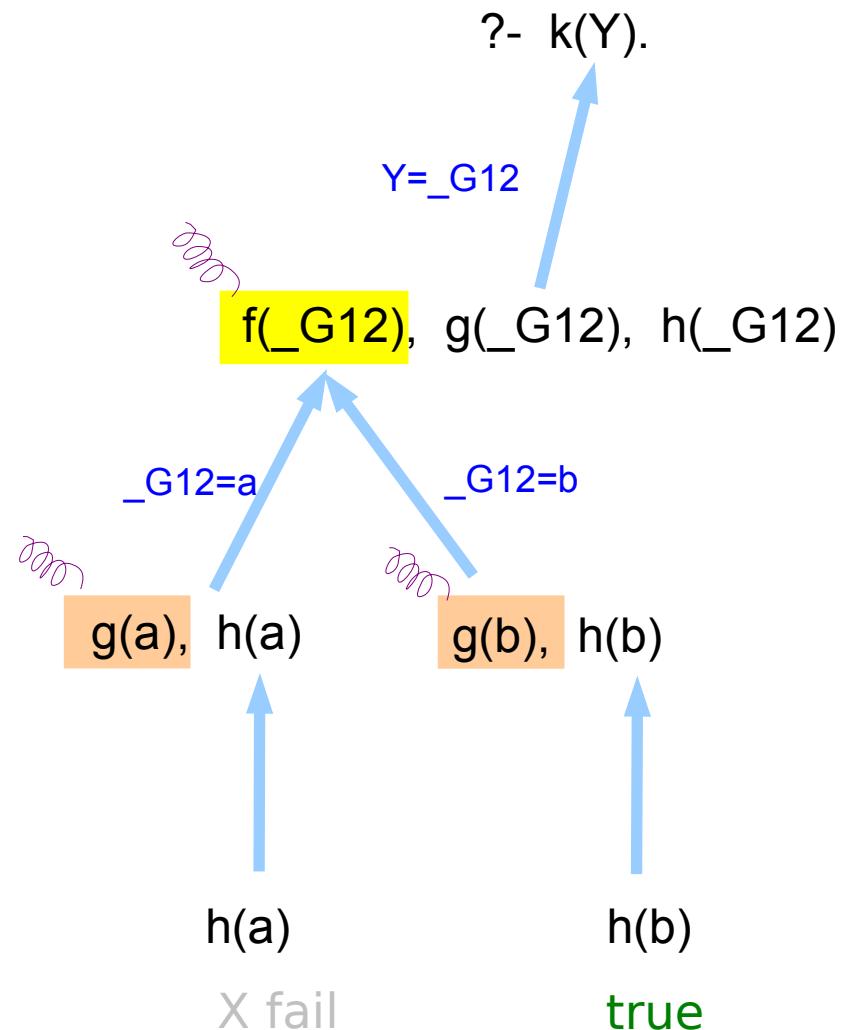
the logical variables in the body have been
bound as a result of the **unification**
Prolog keeps track of **unifying substitutions**

- depth first traversal
- backtracking

Shared Variables

```
f(a).  
f(b).  
  
g(a).  
g(b).  
  
h(b).  
  
k(X) :- f(X), g(X), h(X).
```

When Prolog **unifies** the variable in a **query** to a variable in a **fact** or **rule**, it generates a brand **new variable** (say **_G12**) to represent **the shared variables**.

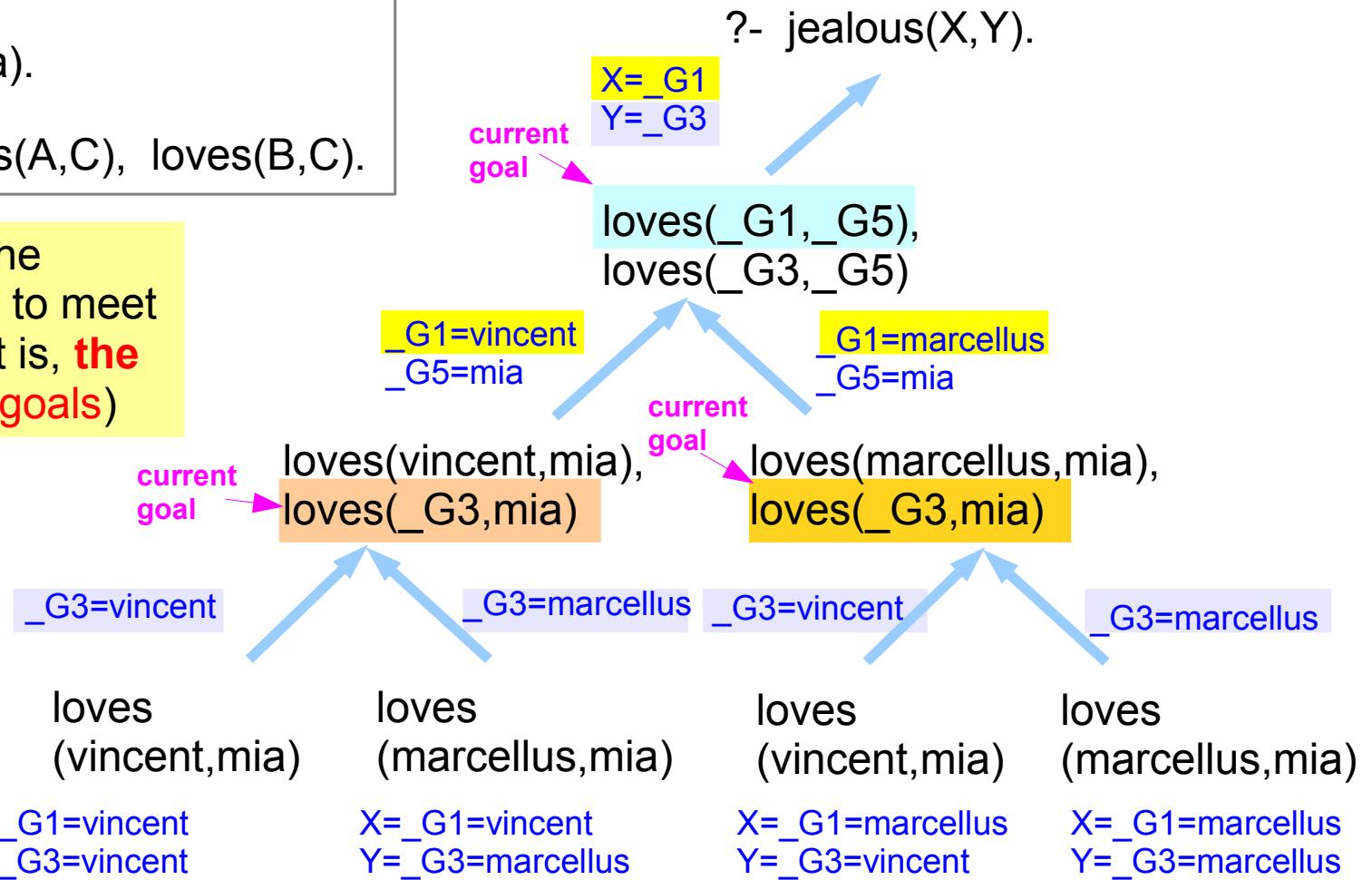


Current Goals

```
loves(vincent,mia).
loves(marcellus,mia).
```

```
jealous(A,B):- loves(A,C), loves(B,C).
```

the **edges** represent the variable instantiations to meet the **current goal** (that is, **the first one in the list of goals**)



Writing Recursive Rules

```
child(bridget,caroline).  
child(caroline,donna).
```

~~descend(X,Y) :- child(X,Y).~~

~~descend(X,Y) :- child(X,Z), child(Z,Y).~~

a non-recursive rule

```
child(anne,bridget).  
child(brIDGET,caroline).  
child(caroline,donna).  
child(donna,emily).
```

~~descend(X,Y) :- child(X,Z_1),
child(Z_1,Z_2),
child(Z_2,Y).~~

~~descend(X,Y) :- child(X,Z_1),
child(Z_1,Z_2),
child(Z_2,Z_3),
child(Z_3,Y).~~

a recursive rule

```
child(anne,brIDGET).  
child(brIDGET,caroline).  
child(caroline,donna).  
child(donna,emily).
```

descend(X,Y) :- child(X,Y).
descend(X,Y) :- child(X,Z), descend(Z,Y).

a problem

a smaller problem

a non-recursive rule

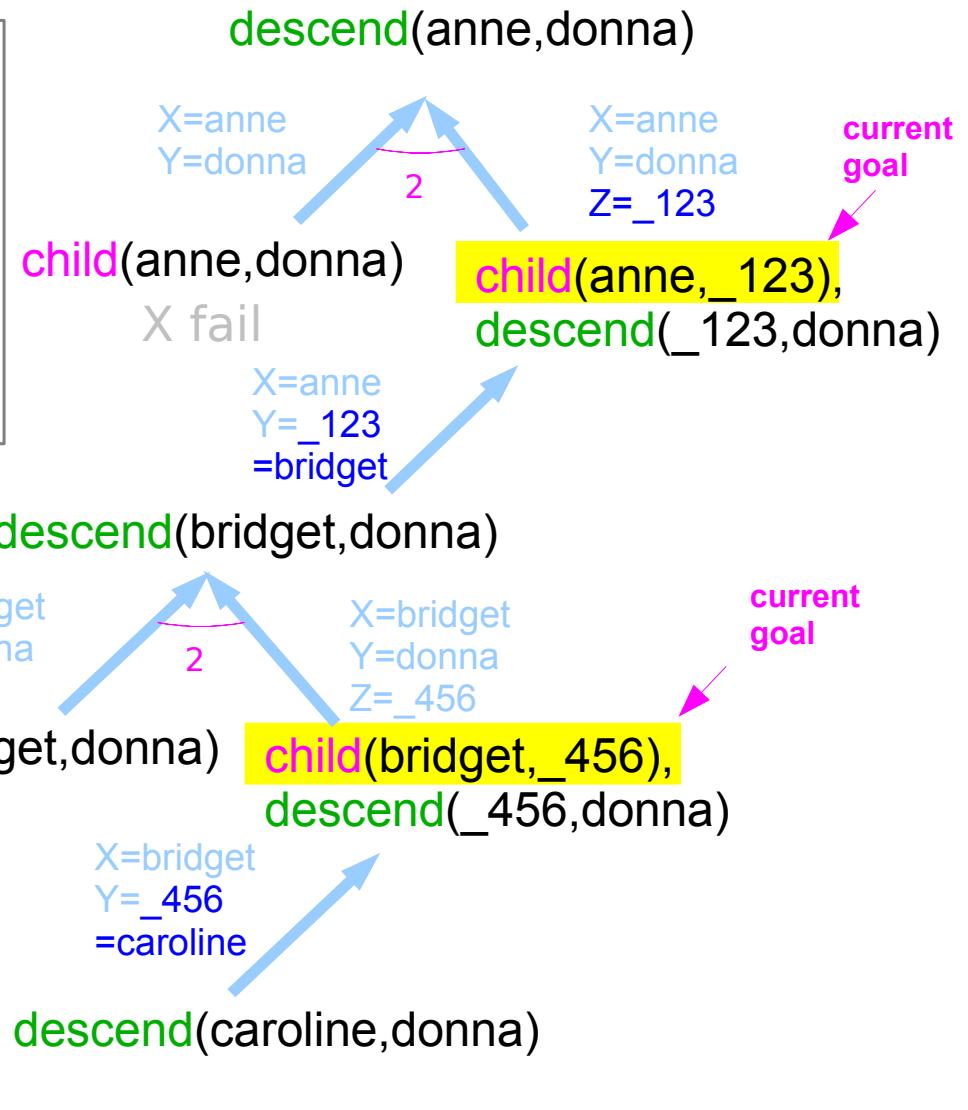
The first rule first

```
child(anne,bridget).
child(brIDGET,caroline).
child(caroline,donna).
child(donna,emily).
```

$\{$

```
descend(X,Y) :- child(X,Y).
descend(X,Y) :- child(X,Z), descend(Z,Y).
```

child(caroline,donna)



A Number System

The number representation using only 4 symbols: 0, succ, (,)

succ(X) : the number X + 1

number 2

number 2

number 4

?- add(succ(succ(0)), succ(succ(0)), succ(succ(succ(succ(0))))).

Yes

number 2

number 1

?- add(succ(succ(0)), succ(0), Y).

Y = succ(succ(succ(0)))

add(0, Y, Y).

add(succ(X), Y, succ(Z)) :- add(X, Y, Z).

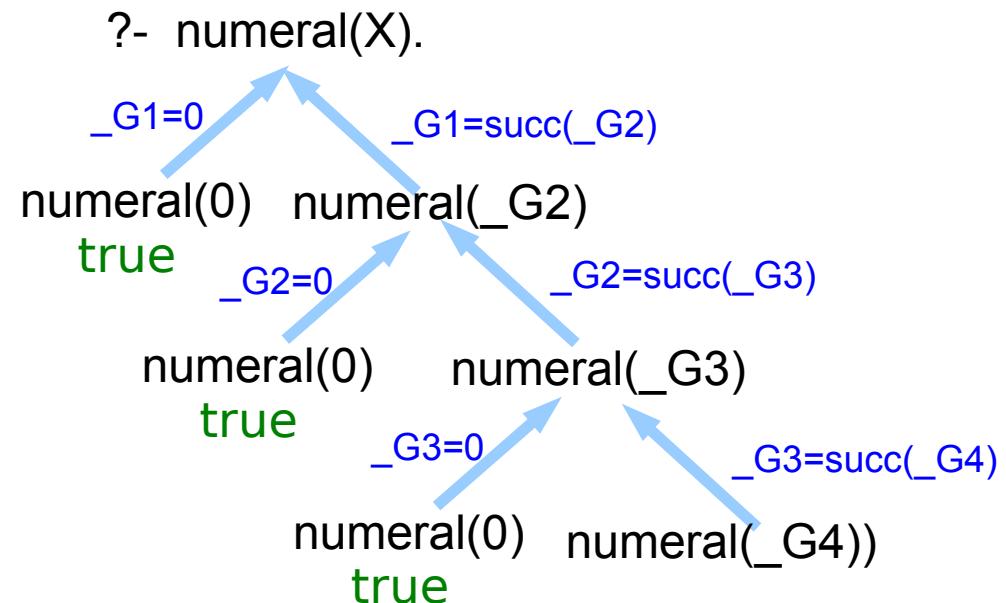
$$0 + Y = Y$$

$$X+1 + Y = Z+1 \quad \leftarrow \quad X + Y = Z$$

Alternative Solutions

```
numeral(0).
```

```
numeral(succ(X)) :-  
    numeral(X).
```



```
X = 0 ;  
X = succ(0) ;  
X = succ(succ(0)) ;  
X = succ(succ(succ(0))) ;  
X = succ(succ(succ(succ(0)))) ;
```

```
numeral(X).
```

```
?- numeral(_G1).
```

```
numeral(succ(_G2)) :- numeral(_G2).
```

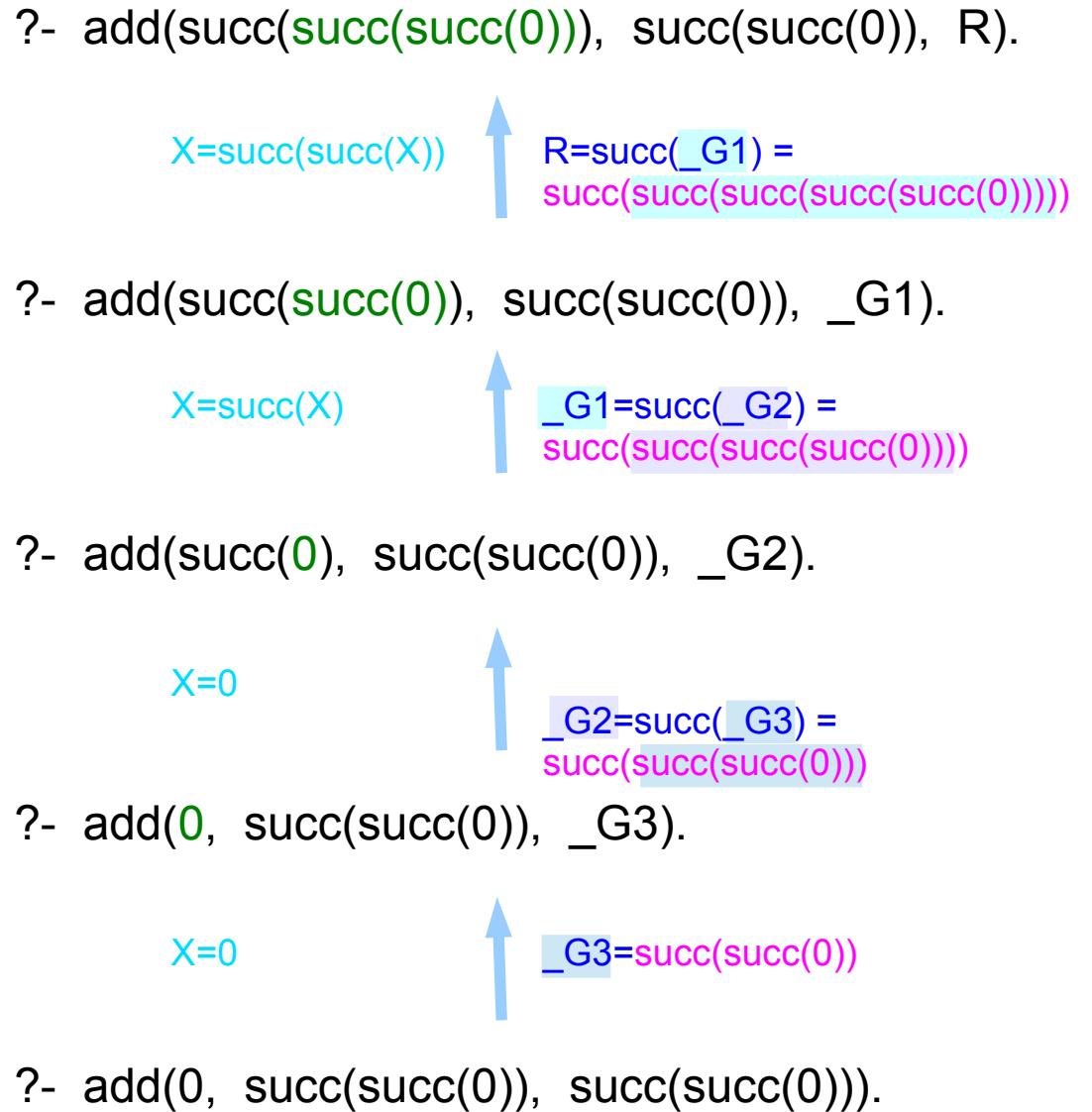
```
numeral(succ(X)) :- numeral(X).
```

Instantiating with a complex term

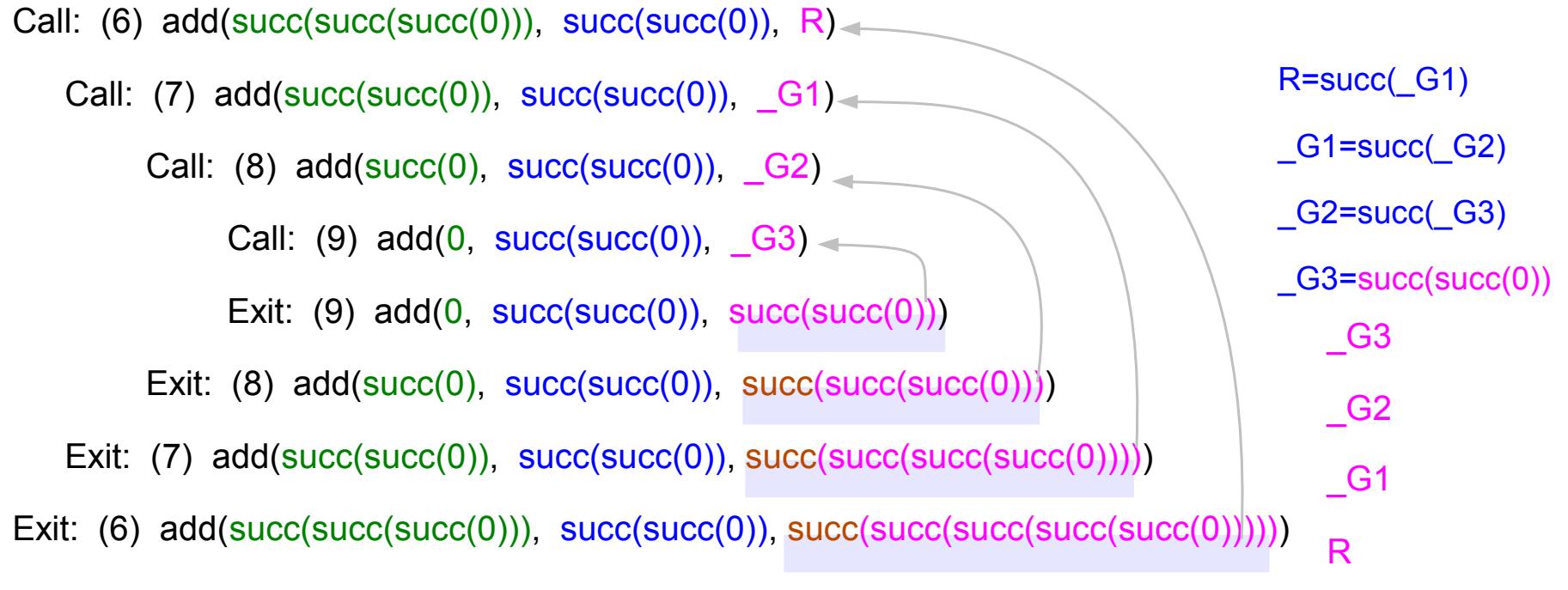
```
add(0, Y, Y).  
add(succ(X), Y, succ(Z))  
:- add(X, Y, Z).
```

R was instantiated to `succ(_G1)`

But that means that R is not a completely uninstantiated variable anymore. It is now a **complex term**, that has **a (uninstantiated) variable** as its argument.



Instantiations On Exiting



`R=succ(_G1)
= succ(succ(_G2))
= succ(succ(succ(_G3)))
= succ(succ(succ(succ(succ(0))))))`

Recursive Definition and Base Case

?- append([a,b,c], [1,2,3], [a,b,c,1,2,3]).

yes

?- append([a,[foo,gibble],c], [1,2,[],b]), [a,[foo,gibble],c,1,2,[],b]).

yes

?- append([a,b,c], [1,2,3], L3).

L3 = [a,b,c,1,2,3]

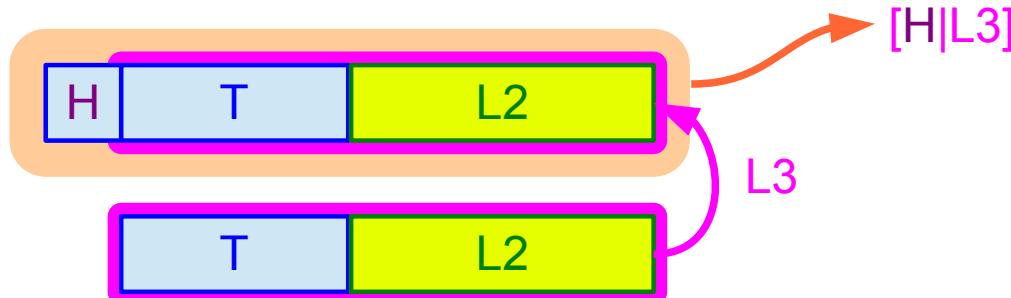
yes

append([], L, L).

append([H|T], L2, [H|L3]) :- append(T, L2, L3).

base case

recursive definition



Instantiation during calling and exiting

```
append([], L, L).
```

```
append([H|T], L2, [H|L3]) :- append(T, L2, L3).
```



[H|T] is instantiated
when the call is made



[H|L3] is instantiated
during exiting the call

```
?- append([a,b], [1,2,3,4], X) .
```

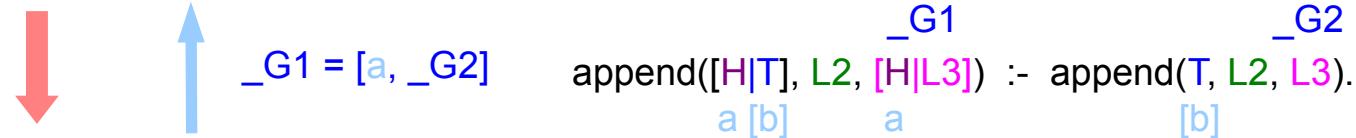
```
append([a, b], [1, 2, 3, 4], _G12)
append([b], [1, 2, 3, 4], _G34)
append([], [1, 2, 3, 4], _G56)
append([], [1, 2, 3, 4], [1, 2, 3, 4])
append([b], [1, 2, 3, 4], [b, 1, 2, 3, 4])
append([a, b], [1, 2, 3, 4], [a, b, 1, 2, 3, 4])
```

```
X = [a, b, c, 1, 2, 3, 4]
yes
```

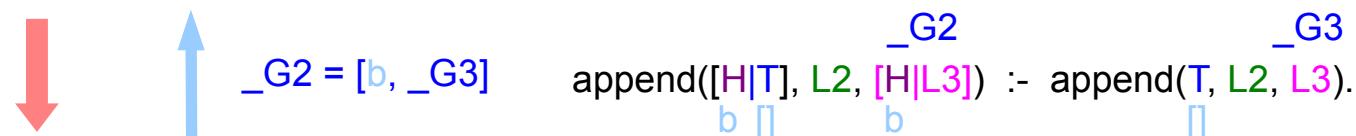
$$\begin{aligned}_G1 &= [a] _G1 \\ &= [a] [b] _G2 \\ &= [a] [b] [c] _G3\end{aligned}$$

Finding subgoals

Goal 1 $\text{append}([a,b], [1,2,3,4], _G1) .$



Goal 2 $\text{append}([b], [1,2,3,4], _G2) .$



Goal 3 $\text{append}([], [1,2,3,4], _G3) .$

$\text{append}([], \text{L}, \text{L}).$

$\text{append}([\text{H}|\text{T}], \text{L2}, [\text{H}|\text{L3}]) :- \text{append}(\text{T}, \text{L2}, \text{L3}).$

\downarrow
[H|T] is instantiated
when the call is made

\uparrow
[H|L3] is instantiated
during exiting the call

Answering subgoals

Goal 1 $\text{append}([a,b], [1,2,3,4], _G1) .$



$_G1 = [a, _G2]$



$_G1 = [a, b, 1,2,3,4]$

Goal 2 $\text{append}([b], [1,2,3,4], _G2) .$



$_G2 = [b, _G3]$



$_G2 = [b, 1,2,3,4]$

Goal 3 $\text{append}([], [1,2,3,4], _G3) .$



$_G3 = [1,2,3,4]$



$\text{append}([], L, L).$

base $\text{append}([], [1,2,3,4], [1,2,3,4]) .$

$\text{append}([], L, L).$

$\text{append}([H|T], L2, [H|L3]) :- \text{append}(T, L2, L3).$



[H|T] is instantiated
when the call is made



[H|L3] is instantiated
during exiting the call

References

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- [2] en.wiktionary.org
- [3] U. Endriss, "Lecture Notes : Introduction to Prolog Programming"
- [4] <http://www.learnprolognow.org/> Learn Prolog Now!
- [5] http://www.csupomona.edu/~jrfisher/www/prolog_tutorial
- [6] www.cse.unsw.edu.au/~billw/cs9414/notes/prolog/intro.html
- [7] www.cse.unsw.edu.au/~billw/dictionaries/prolog/negation.html