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**graph traversal (graph search)** refers to the process of <u>visiting</u> (checking and/or updating) each **vertex** in a graph.

Such traversals are <u>classified</u> by the <u>order</u> in which the vertices are visited.

Tree traversal is a special case of graph traversal.

A depth-first search (**DFS**) is an algorithm for traversing a finite graph.

DFS visits the **child vertices** <u>before</u> visiting the **sibling vertices**;

that is, it traverses the **depth** of any particular path <u>before</u> exploring its **breadth**.

A **stack** (often the program's call stack via recursion) is generally used when implementing the algorithm.

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The algorithm begins with a chosen "**root**" vertex;

it then iteratively transitions from the **current** vertex to an **adjacent**, **unvisited** vertex, until it can <u>no longer</u> find an **unexplored vertex** to transition to from its current location.

The algorithm then **backtracks** along previously **visited vertices**, until it finds a vertex connected to yet more uncharted territory.

It will then proceed down the **new path** as it had before, **backtracking** as it encounters **dead-ends**, and ending only when the algorithm has backtracked past the original "root" vertex from the very first step.

A breadth-first search (**BFS**) is another technique for traversing a finite graph.

BFS visits the **neighbor** vertices <u>before</u> visiting the **child** vertices

a queue is used in the search process

This algorithm is often used to find the **shortest path** from one vertex to another.

### **Depth First Search Example**



https://en.wikipedia.org/wiki/Graph\_traversal



#### **Breadth First Search Example**



https://en.wikipedia.org/wiki/Graph\_traversal



### General Graph Search Algorithm – 1

```
Search(Start, isGoal, Criteria)
insert(Start, Open);
repeat
if (empty(Open)) then return fail;
select node from Open using Criteria;
mark node as <u>visited</u>;
if (isGoal(node)) then return node;
for each child of node do
if (child <u>not</u> already <u>visited</u>)
then insert(child, Open);
```

https://courses.cs.washington.edu/courses/cse326/08wi/a/lectures/lecture13.pdf

https://en.wikiversity.org/wiki/Artificial\_intelligence/Lecture\_aid



https://en.wikiversity.org/wiki/Artificial\_intelligence/Lecture\_aid



### **Possible duplication**

https://en.wikiversity.org/wiki/Artificial\_intelligence/Lecture\_aid



### Must check before expansion

https://en.wikiversity.org/wiki/Artificial\_intelligence/Lecture\_aid



### General Graph Search Algorithm – 1

```
Search(Start, isGoal, Criteria)
insert(Start, Open);
repeat
if (empty(Open)) then return fail;
select node from Open using Criteria;
mark node as <u>visited</u>;
if (isGoal(node)) then return node;
for each child of node do
if (child <u>not</u> already <u>visited</u>)
then insert(child, Open);
```

Remedy 1: check if visited when selecting

Remedy 2: check redundant nodes

https://courses.cs.washington.edu/courses/cse326/08wi/a/lectures/lecture13.pdf

# DFS-1 (Depth First Search)

Open list – use a stack Select with Criteria – **pop** 

<b>DFS</b> (Start, isGoal)	
<b>push</b> (Start, <mark>Open</mark> );	// push
repeat	
if ( <b>empty(Open</b> )) <b>then return</b> fail;	
node := <b>pop</b> (Open);	// рор
mark node as <u>visited;</u>	
if ( <b>isGoal</b> (node)) then return node;	
for each child of node do	
if (child <u>not</u> already <u>visited</u> ) then	
<pre>push(child, Open);</pre>	// push

https://courses.cs.washington.edu/courses/cse326/08wi/a/lectures/lecture13.pdf

# DFS-1 Example (1)



https://en.wikipedia.org/wiki/Graph\_traversal

# DFS-1 Example (2)



https://en.wikipedia.org/wiki/Graph\_traversal

### BFS-1 (Breadth First Search)

Open list – use a FIFO Select with Criteria – **dequeue** 

```
BFS(Start, isGoal)
enqueue(Start, Open); // enqueue
repeat
if (empty(Open)) then return fail;
node := dequeue(Open); // dequeue
mark node as visited;
if (isGoal(node)) then return node;
for each child of node do
    if (child not already visited) then
    enqueue(child, Open); // enqueue
```

https://courses.cs.washington.edu/courses/cse326/08wi/a/lectures/lecture13.pdf

# BFS-1 Example (1)



https://en.wikipedia.org/wiki/Graph\_traversal

# BFS-1 Example (2)



https://en.wikipedia.org/wiki/Graph\_traversal

#### Graph Search (6A)



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### General Graph Search Algorithm – 2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    LIST = {s}
while LIST ≠ Ø do
    select a node i in LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        add node j to the end of LIST;
    else
        delete node i from LIST
```



**BFS : select** the first node i in LIST;



### Admissible arc

pred(j) is a node that precedes j on some path from s;

A node is either **marked** or **unmarked**.

Initially only node s is marked.

If a node is marked, it is reachable from node s.

An arc  $(i,j) \in A$  is **admissible** if node i is <u>marked</u> and j is <u>not</u>.



# LIST

Before a node is <u>added</u> into LIST, the node is **marked** 

LIST contains only the marked nodes

thus, the <u>selected</u> node **i** is **marked** already

The node **j** incident to the **admissible** arc(**i**,**j**) must be **unmarked** 

This node **j** is **marked** and <u>added</u> into LIST

In this way, LIST contains only **marked** and **non-repeating** nodes

Check before inserting



### DFS-2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    push s onto LIST
while LIST ≠ Ø do
    pop a node i from LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        push(node j) onto LIST;
    else
        delete node i from LIST
```

# DFS-2 Example (1)



https://en.wikipedia.org/wiki/Graph\_traversal

# DFS-2 Example (2)



https://en.wikipedia.org/wiki/Graph\_traversal



### BFS-2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    enqueue s onto LIST
while LIST ≠ Ø do
    dequeue node i from LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        enqueue node j onto LIST;
    else
        delete node i from LIST
```

# BFS-2 Example (1)



https://en.wikipedia.org/wiki/Graph\_traversal

# BFS-2 Example (2)



https://en.wikipedia.org/wiki/Graph\_traversal

#### Graph Search (6A)



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### **DFS** Pseudocode

1 procedure DFS(G, v):

- 2 label v as explored
- 3 for all edges e in G.incidentEdges(v) do
- if edge e is unexplored then 4
- 5 w - G.adjacentVertex(v, e)
- 6 if vertex w is unexplored then 7
  - label e as a discovered edge
- 8 recursively call DFS(G, w)
- 9 else
- 10 label e as a back edge

### **BFS Pseudocode**

1 procedure BFS(G, v):

- 2 create a queue Q
- 3 enqueue v onto Q
- 4 mark v
- 5 while Q is not empty:
- 6 t ← Q.dequeue()
- 7 if t is what we are looking for:
- 8 return t
- 9 for all edges e in G.adjacentEdges(t) do
- 12  $o \leftarrow G.adjacentVertex(t, e)$
- if o is not marked:
- 14 mark o
- 15 enqueue o onto Q
- 16 return null

#### References

