

Applications of Arrays (1A)

Copyright (c) 2010 - 2020 Young W. Lim.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.2 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled "GNU Free Documentation License".

Please send corrections (or suggestions) to youngwlim@hotmail.com.

This document was produced by using LibreOffice.

2-d Arrays

2-d array definition

```
int c [4][4];
```

A matrix view

	col 0	col 1	col 2	col 3
row 0	c [0][0]	c [0][1]	c [0][2]	c [0][3]
row 1	c [1][0]	c [1][1]	c [1][2]	c [1][3]
row 2	c [2][0]	c [2][1]	c [2][2]	c [2][3]
row 3	c [3][0]	c [3][1]	c [3][2]	c [3][3]

2-d array access via pointers

```
int c [4][4];
```

1. recursive pointers

```
c [ i ][ j ]
```

```
(*(c+i))[ j ]
```

→ `int (*p)[4];`

```
*(c[ i ]+ j)
```

```
*(*(c+i)+ j)
```

→ `int **q;`

```
int *p = c[0] ;
```

2. linear array pointers

```
p[ i*4 + j ]
```

```
*(p+ i*4 + j)
```

1-d array names

```
int    a [4];
```

```
int    c [4] [4];
```



The value of `a` is the starting address of a 4 element array

`a`: pointer to the first element



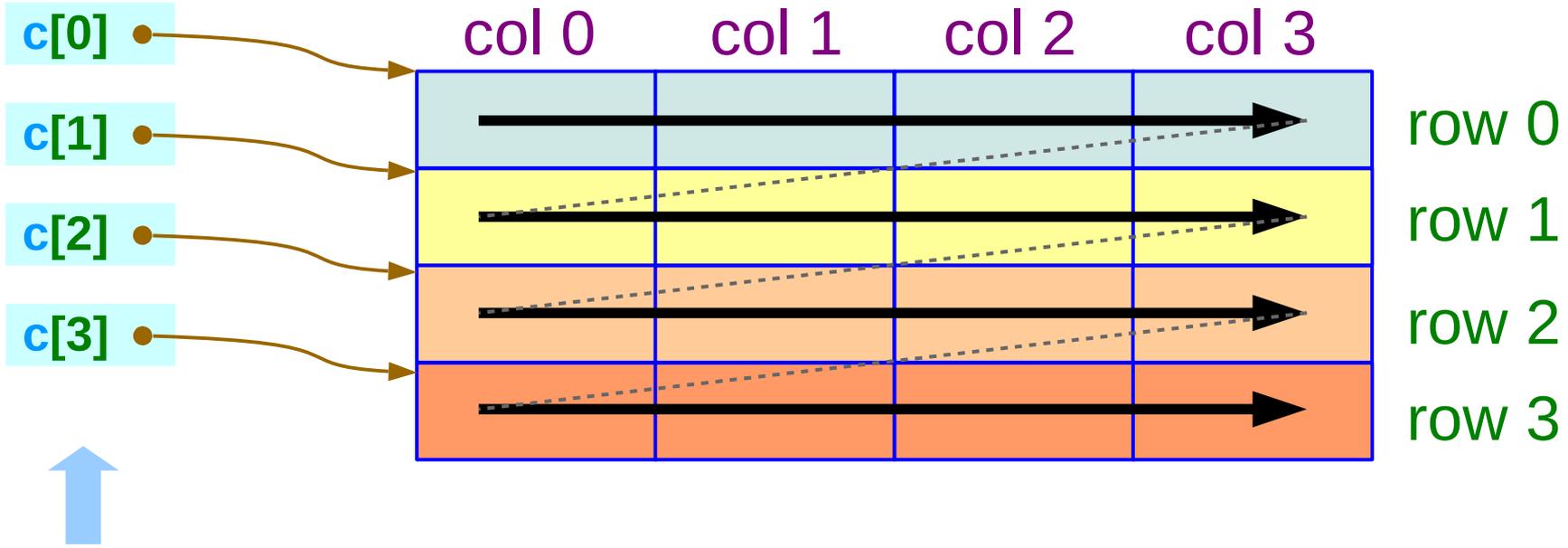
Each value of `c[i]` is the starting address of a 4 element array

`c[i]`: pointer to the first element

2-d array as a matrix

```
int c[4][4];
```

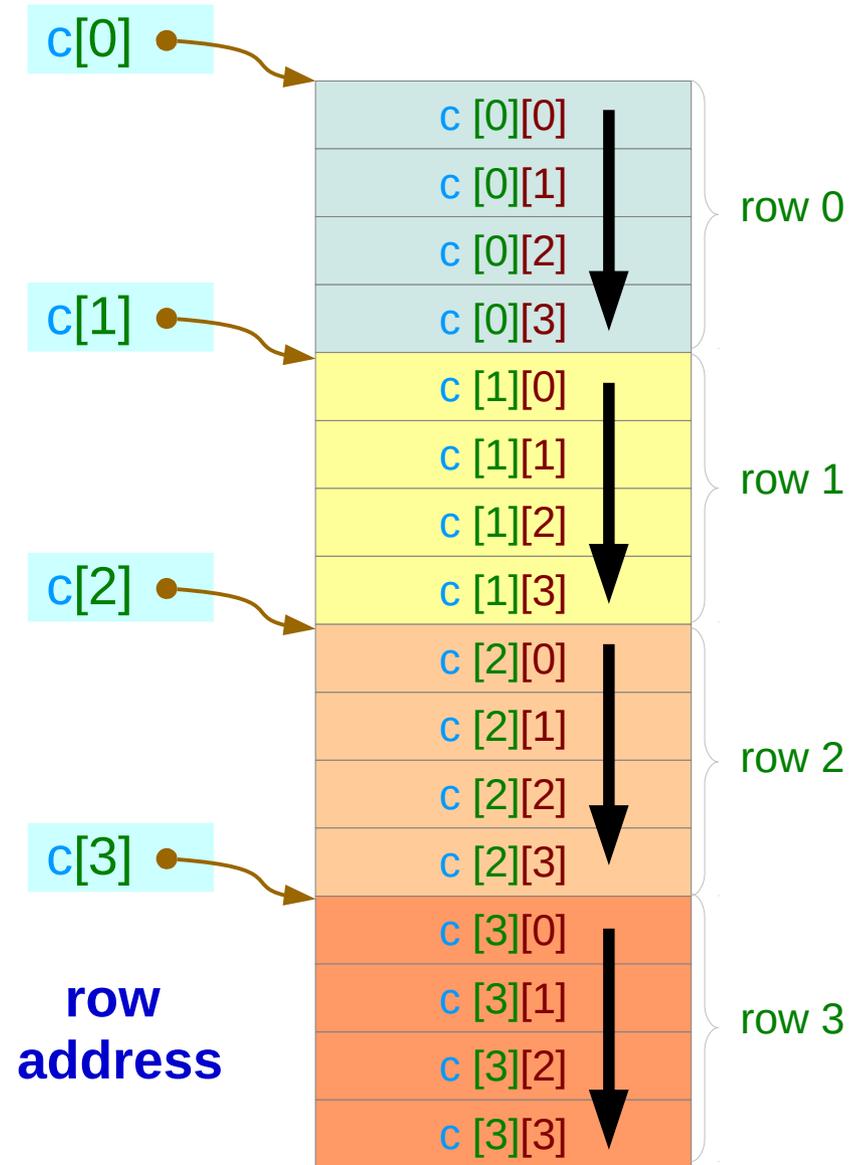
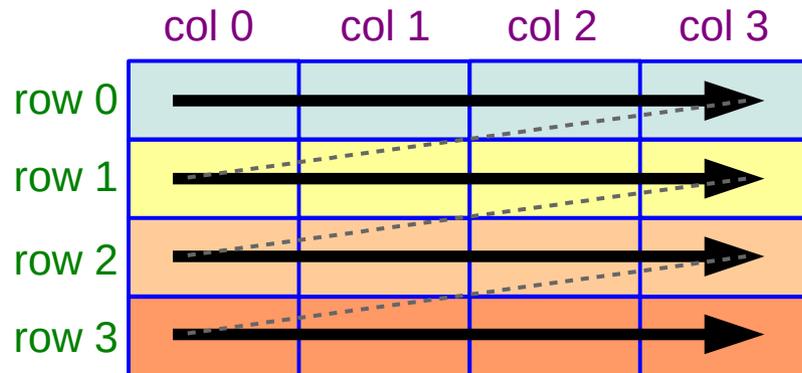
row major ordering



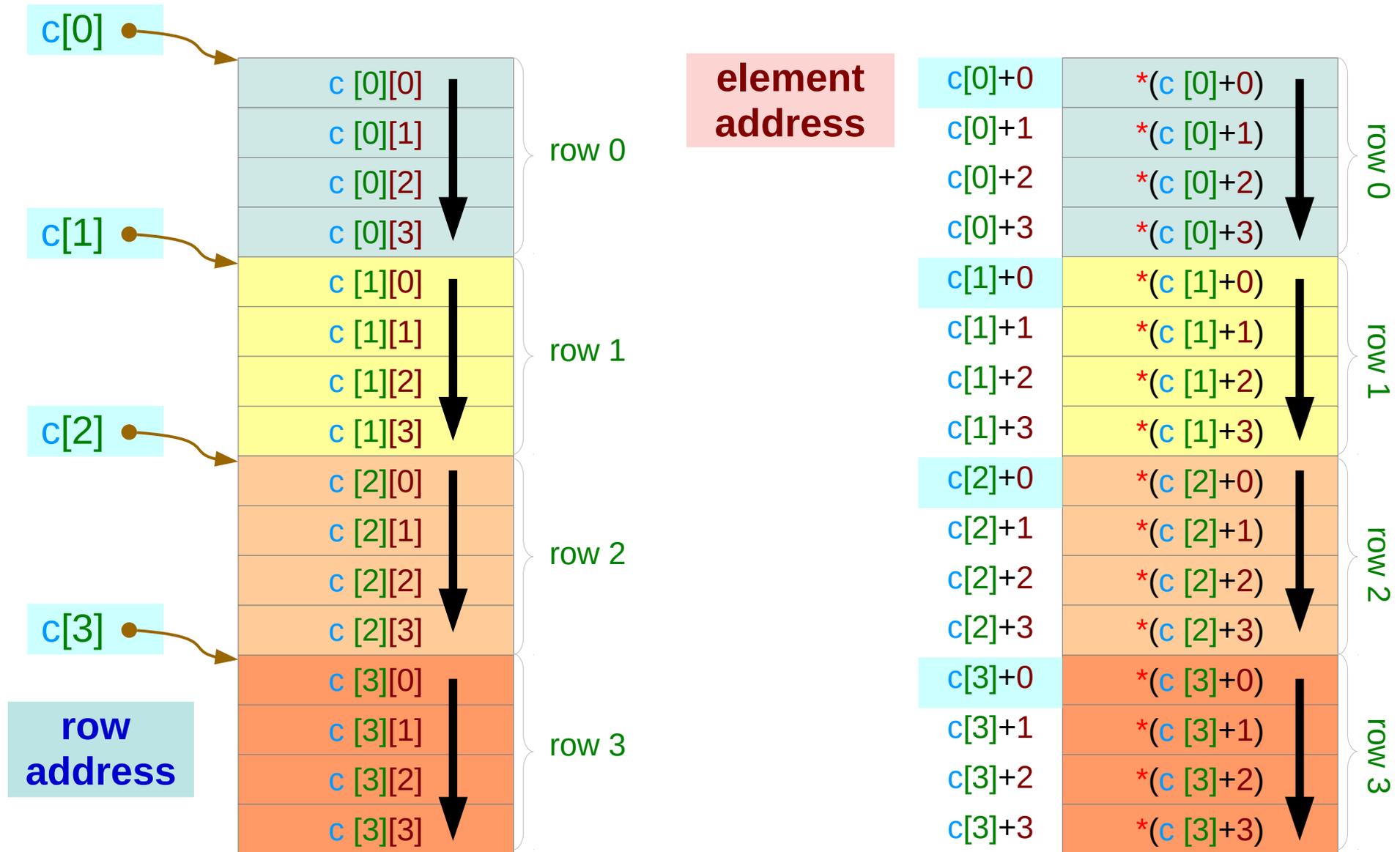
consider each $c[i]$ as a pointer to the first element of each 4 element array

2-d array stored as a linear array

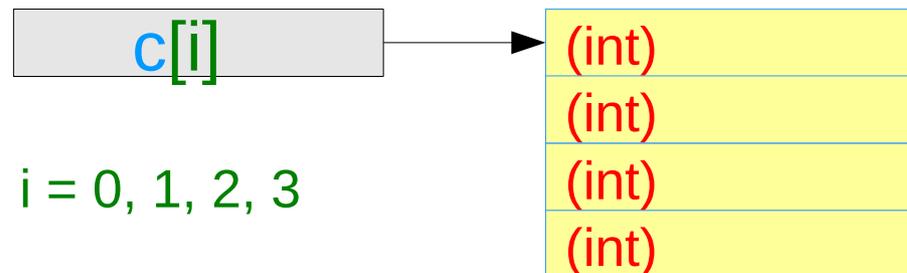
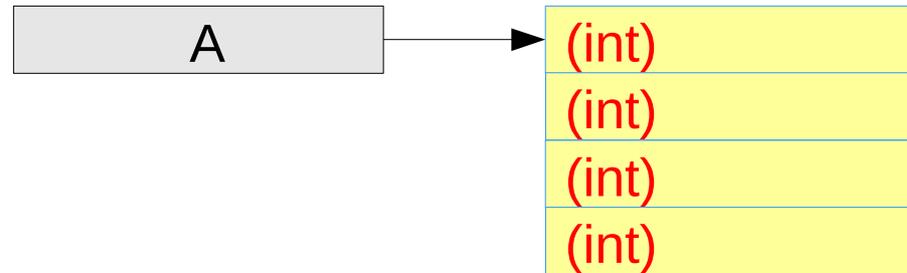
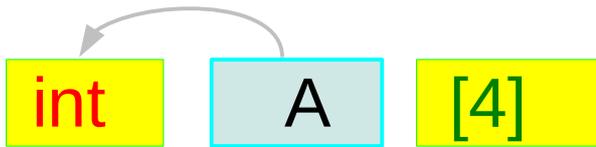
```
int c [4][4];
```



Row Address and Element Address

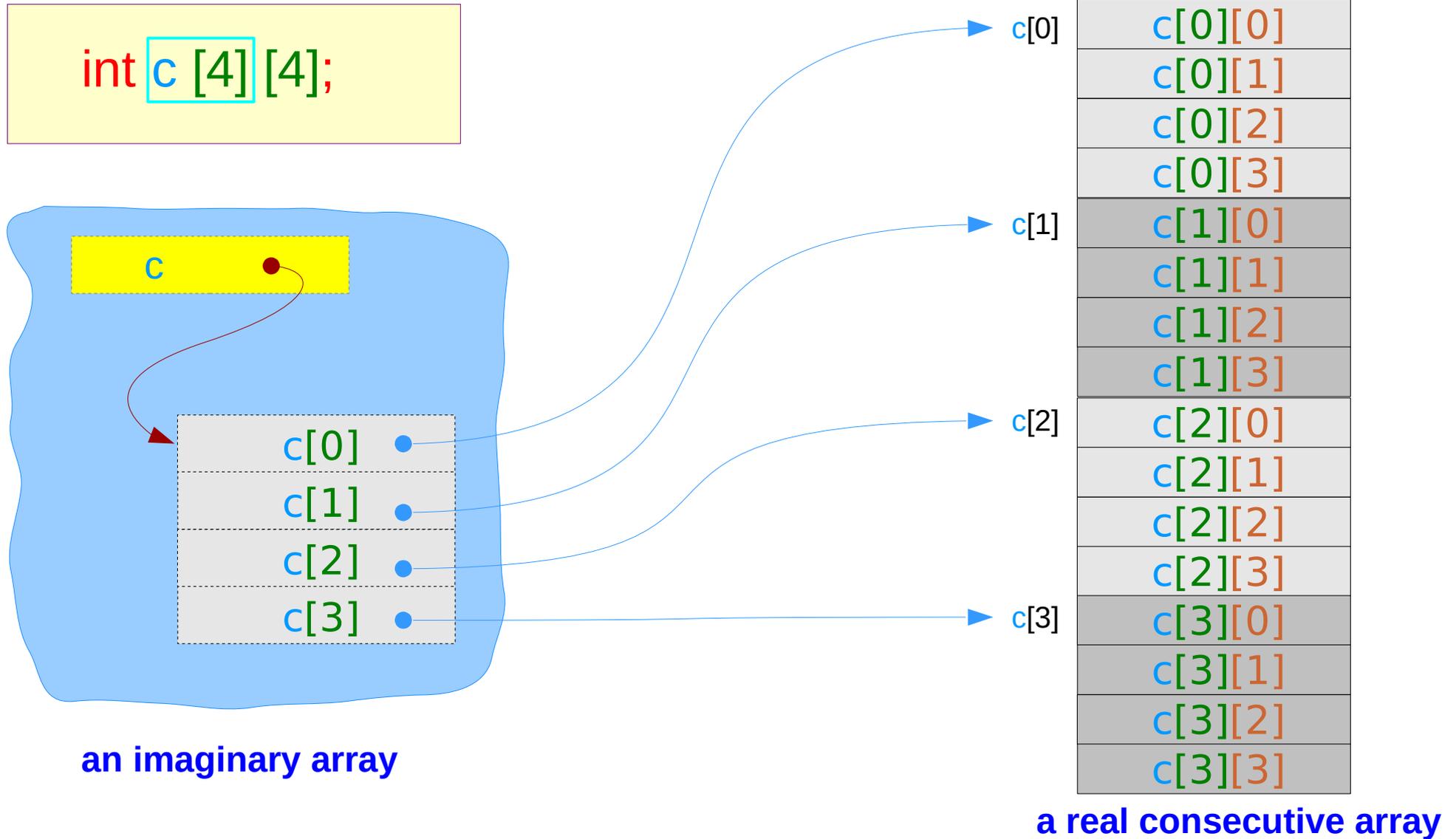


Addresses of 1-d arrays with 4 integer elements

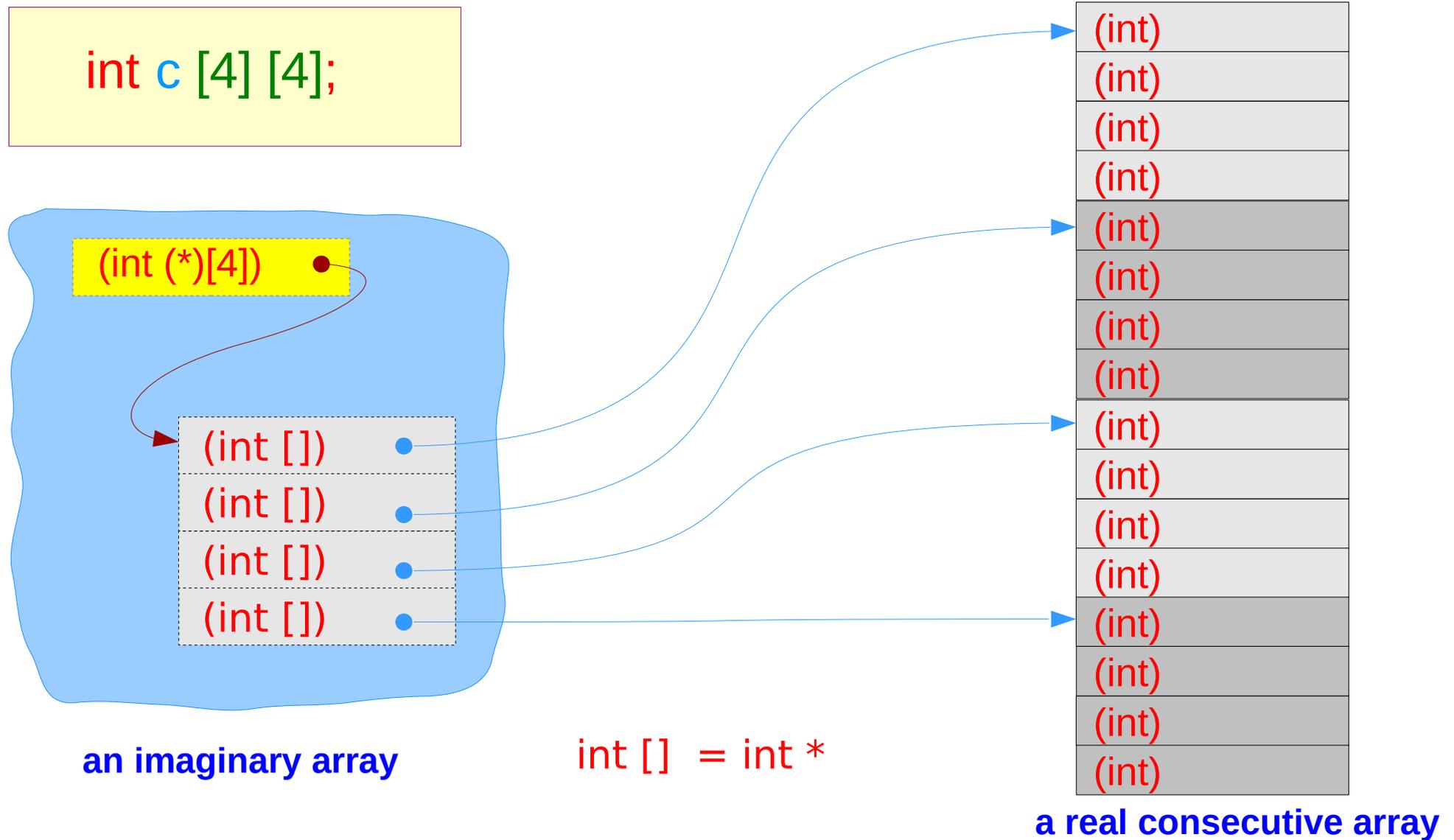


```
int c [4] [4];
```

A 2-d Array – a variable view

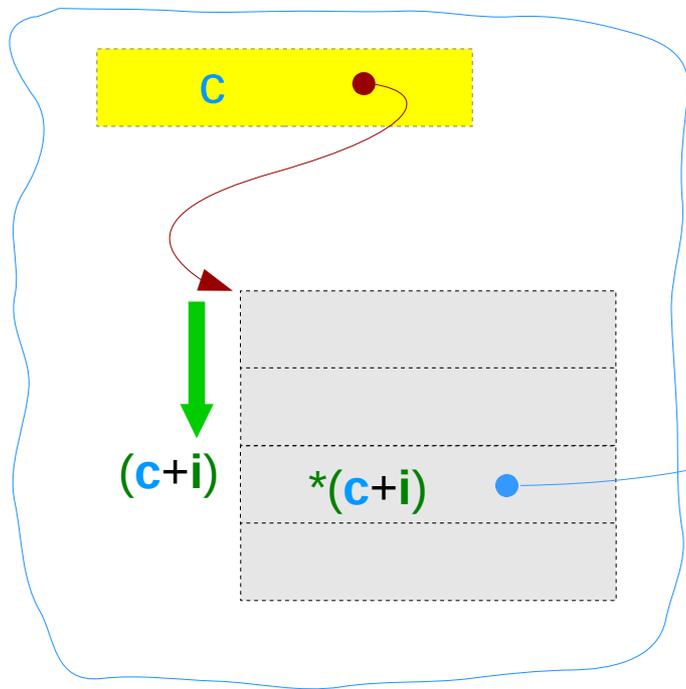


A 2-d Array – a type view

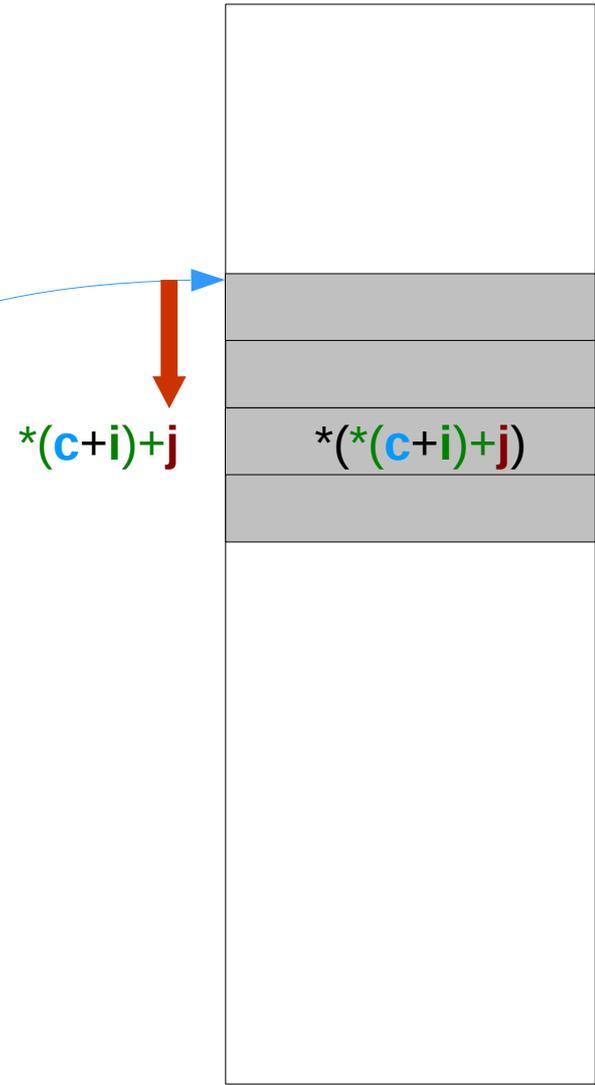


A 2-d Array – an index view

```
int c [4] [4];
```



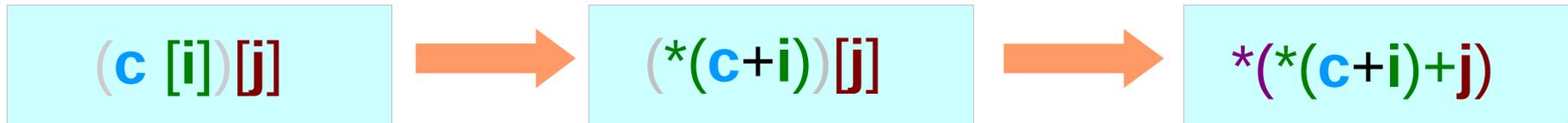
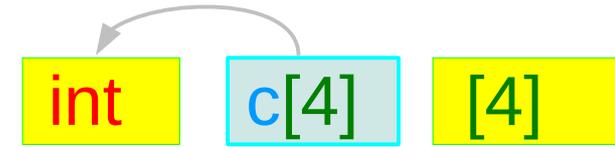
an imaginary array



a real consecutive array

2-d array access via a double indirection

```
int c [4] [4];
```



$$(c [i]) = (*(c+i))$$

contiguous memory
locations are assumed

$$(_) [j] = *((_) + j)$$

contiguous memory
locations are assumed

Recursive Pointer Conversions

```
int c [4][4];
```

```
c[i] = *(c+i) hold row addresses
```

assumption

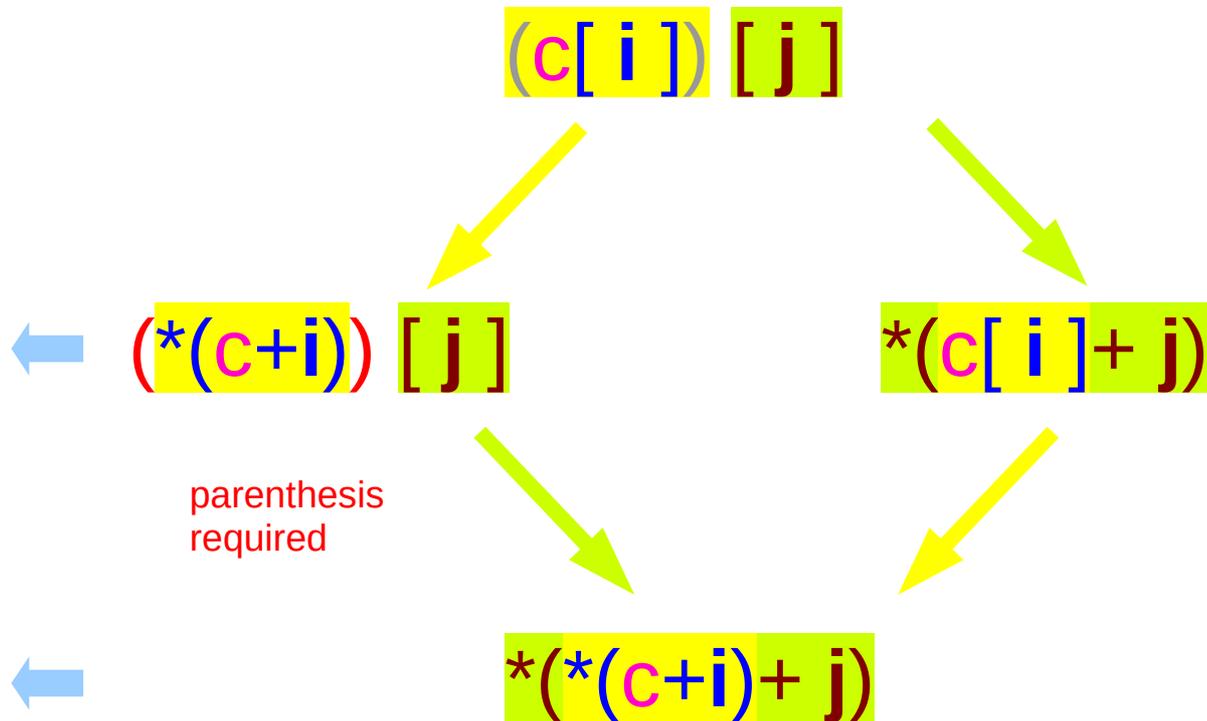
contiguous memory locations are assumed

array pointer

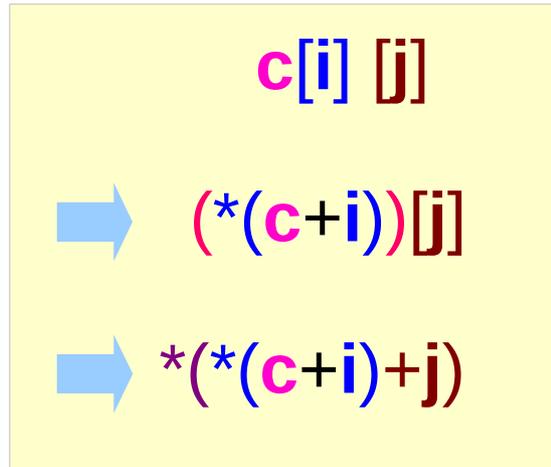
```
int (*p)[4];
```

double pointer

```
int **q;
```



2-d array access via recursive pointers (double indirection)



$*c+i$: row address

$*(*c+i)+j$: element address

first select a row, then a column

$c[0]+0 = *(c+0)+0$ $c[0][0]$

$c[0]+1 = *(c+0)+1$ $c[0][1]$

$c[0]+2 = *(c+0)+2$ $c[0][2]$

$c[0]+3 = *(c+0)+3$ $c[0][3]$

$c[1]+0 = *(c+1)+0$ $c[1][0]$

$c[1]+1 = *(c+1)+1$ $c[1][1]$

$c[1]+2 = *(c+1)+2$ $c[1][2]$

$c[1]+3 = *(c+1)+3$ $c[1][3]$

$c[2]+0 = *(c+2)+0$ $c[2][0]$

$c[2]+1 = *(c+2)+1$ $c[2][1]$

$c[2]+2 = *(c+2)+2$ $c[2][2]$

$c[2]+3 = *(c+2)+3$ $c[2][3]$

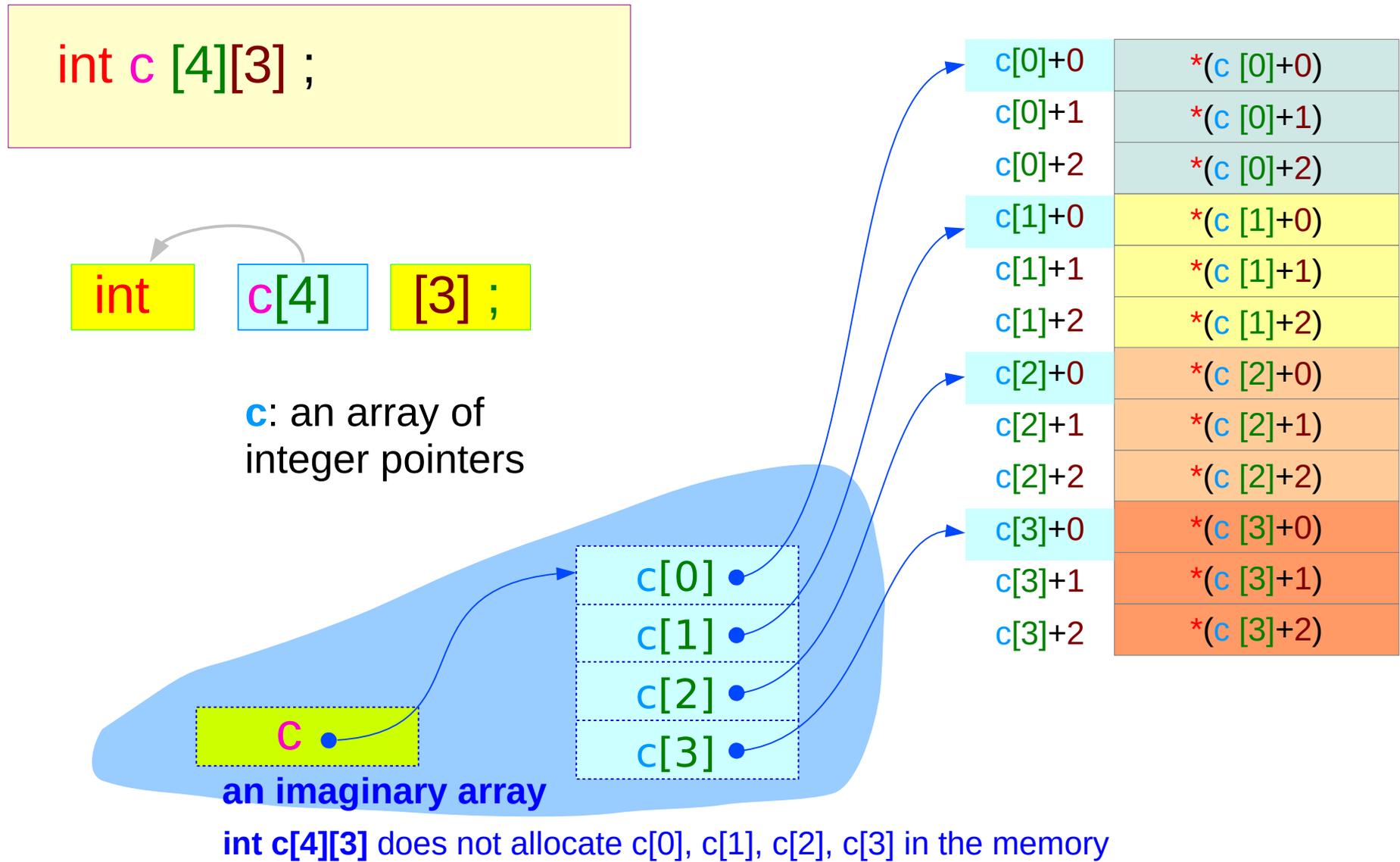
$c[3]+0 = *(c+3)+0$ $c[3][0]$

$c[3]+1 = *(c+3)+1$ $c[3][1]$

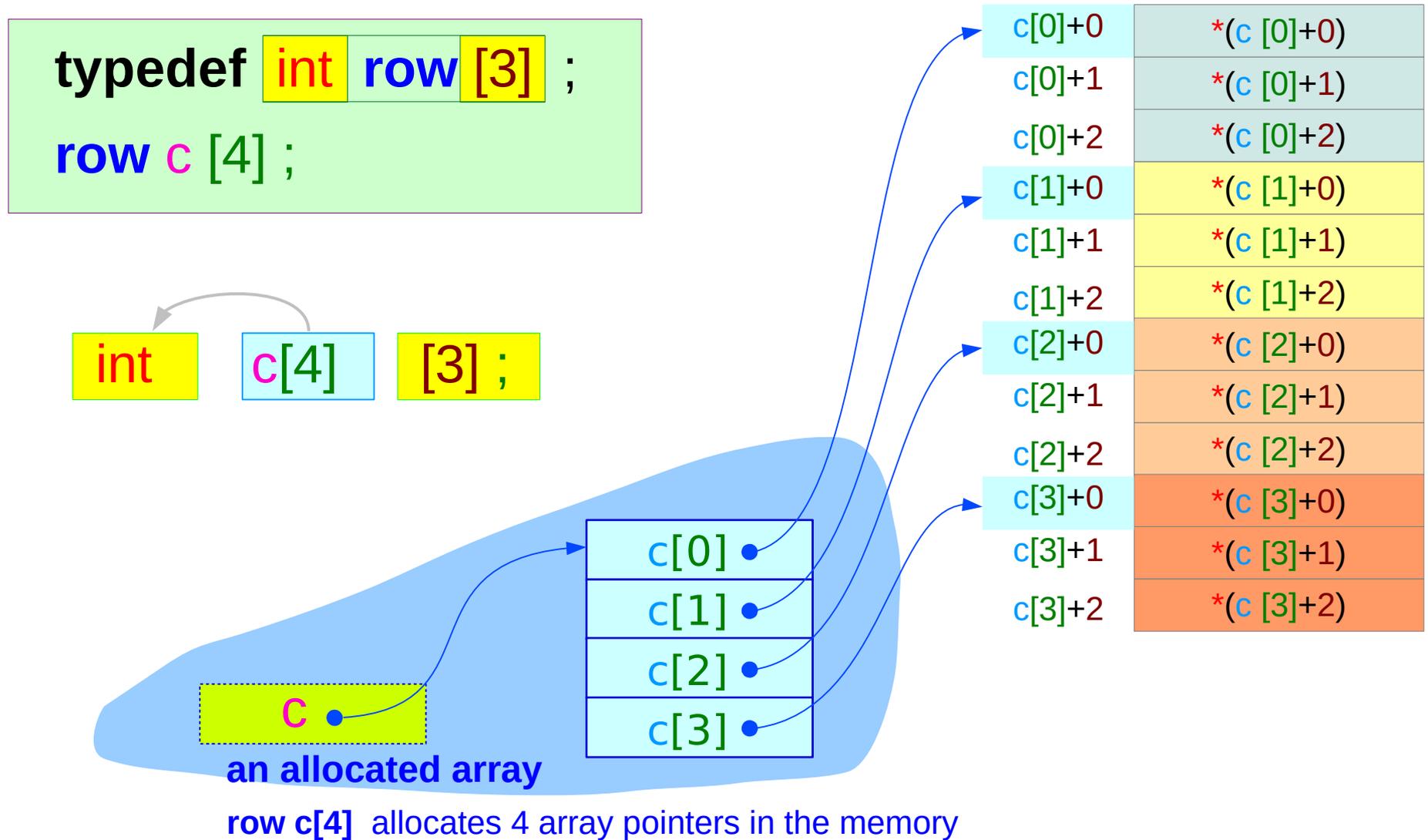
$c[3]+2 = *(c+3)+2$ $c[3][2]$

$c[3]+3 = *(c+3)+3$ $c[3][3]$

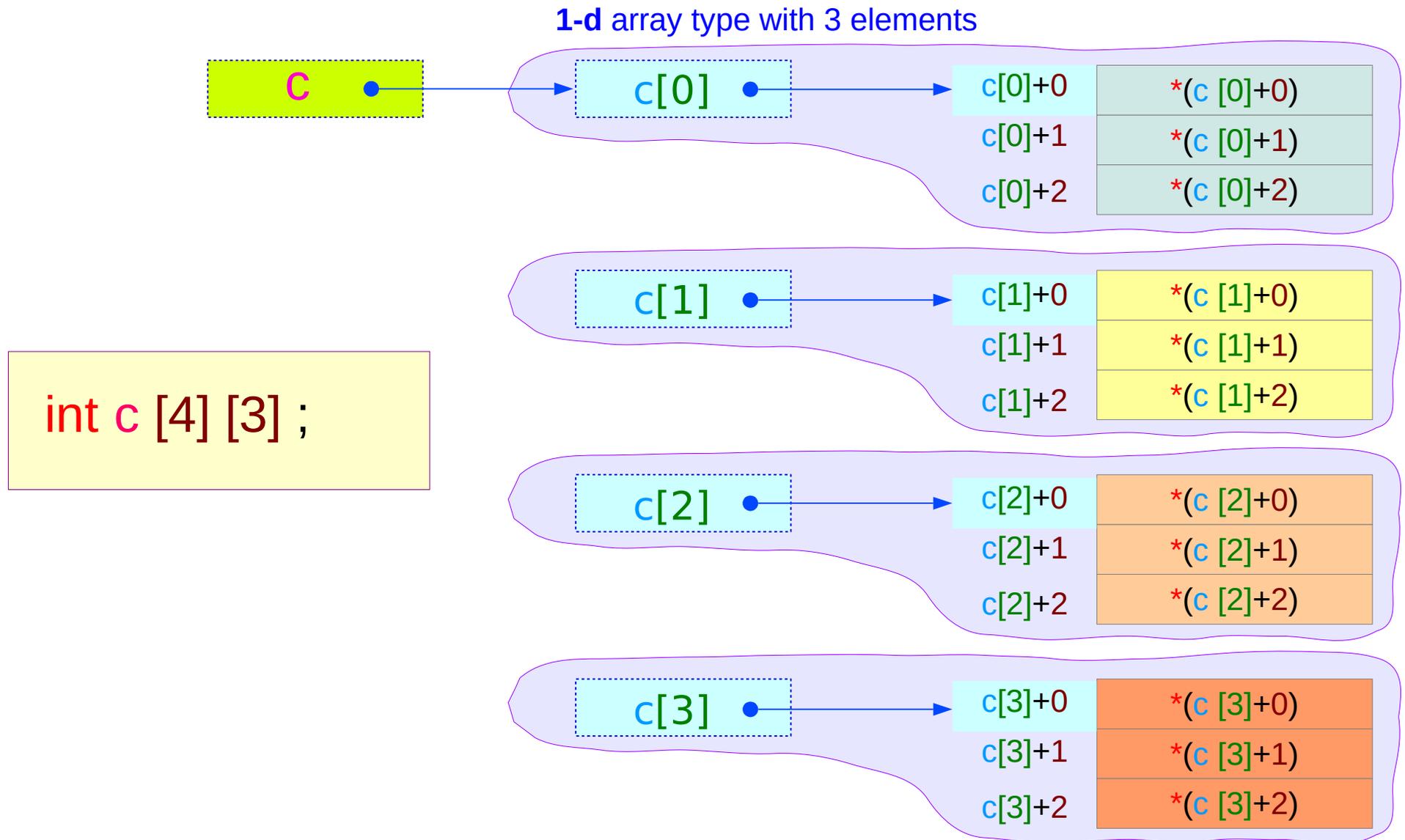
Nested arrays



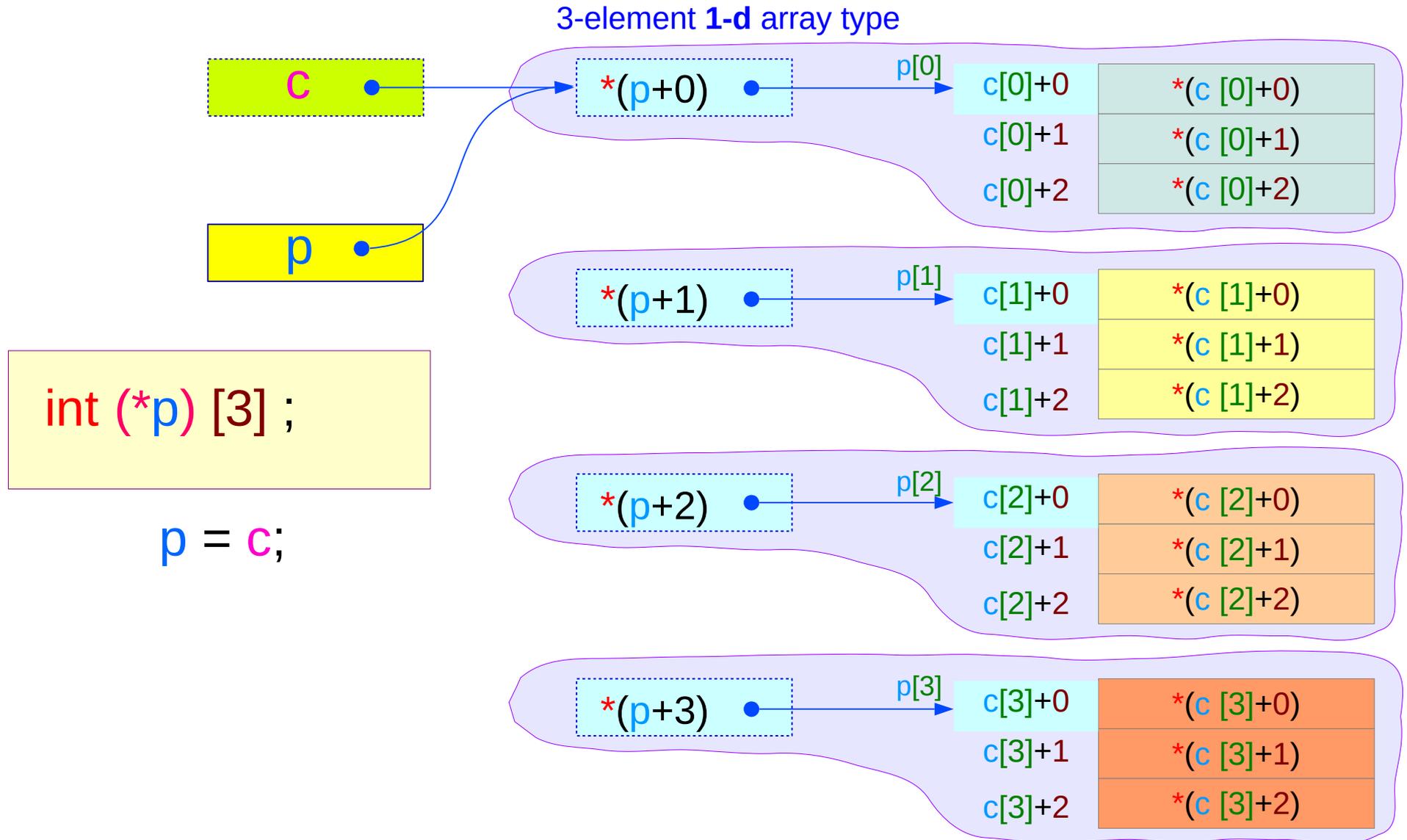
Nested array declared explicitly



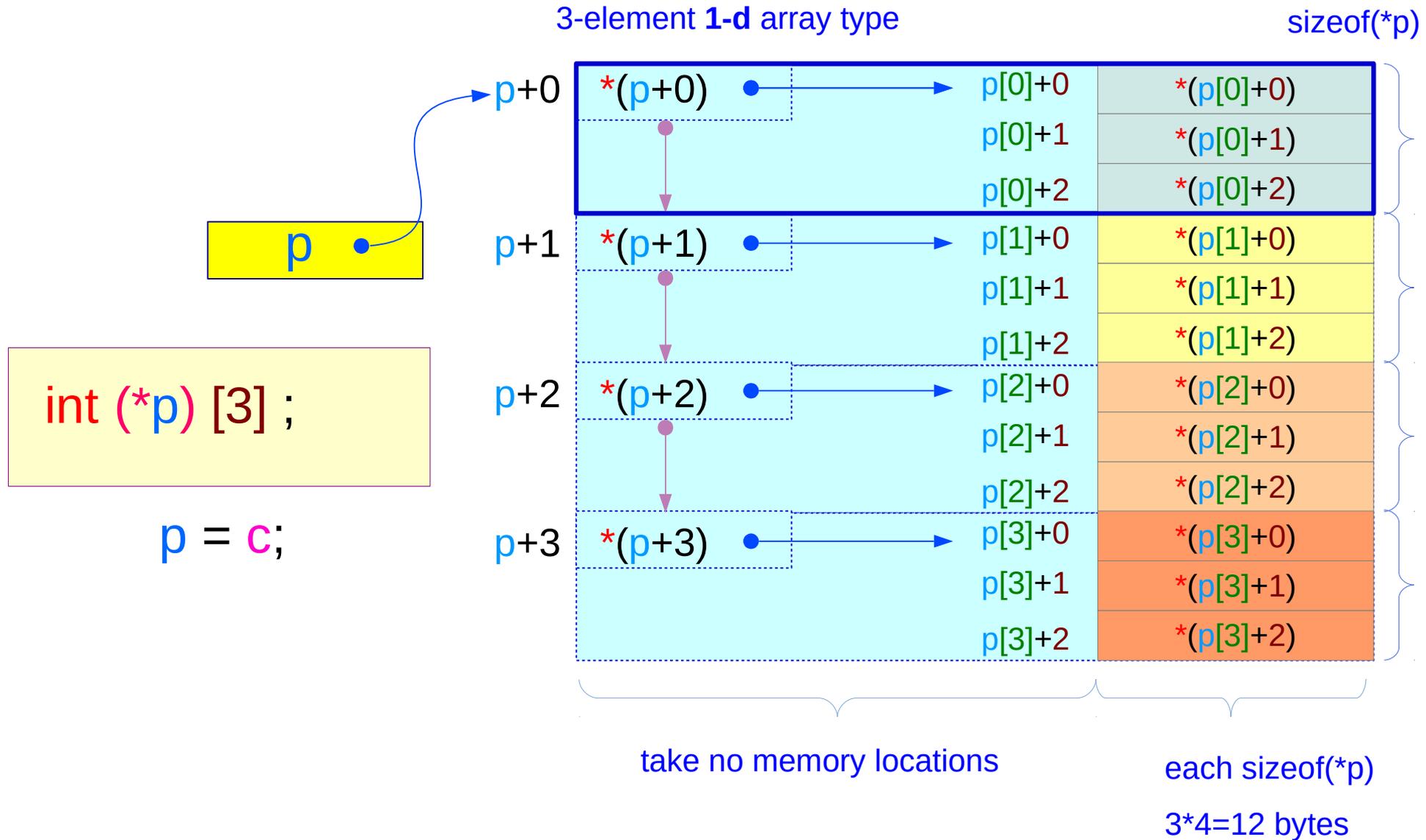
2-d Array : rows of 1-d arrays



Pointer to 1-d arrays



Incrementing the 1-d array pointer p

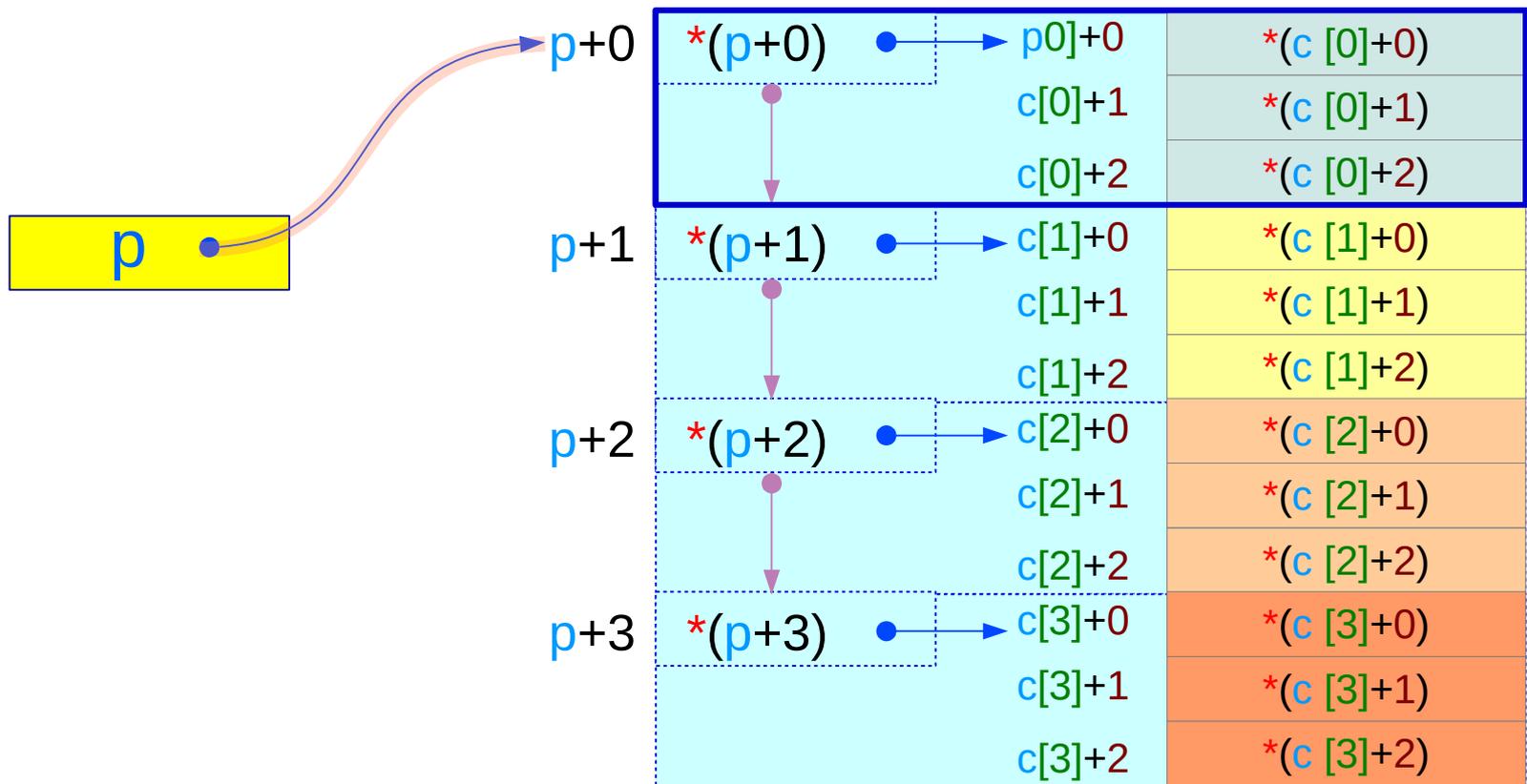


2-d array access via a 1-d array pointer

```
int c [4] [4];
```

```
int (*p) [4];
```

```
p = c;
```



2-d array access via a 1-d array pointer

```
int c [4] [4];
```

```
int (*p) [4];
```

$(c [i]) [j]$



$(*(c+i)) [j]$



$*(*(c+i)+j)$

$p = c$

```
p[0]=c[0],  
p[1]=c[1],  
p[2]=c[2],  
p[3]=c[3];
```

equivalence

$(p [i]) [j]$



$(*(p+i)) [j]$



$*(*(p+i)+j)$

2-d array access via a double pointer

```
int c [4] [4];
```

```
int **p, *q[4];
```

$(c [i])[j]$



$(*(c+i))[j]$



$*(*(c+i)+j)$

$p = q;$

$q[0]=c[0],$
 $q[1]=c[1],$
 $q[2]=c[2],$
 $q[3]=c[3];$

must be allocated
and initialized

$(p [i])[j]$

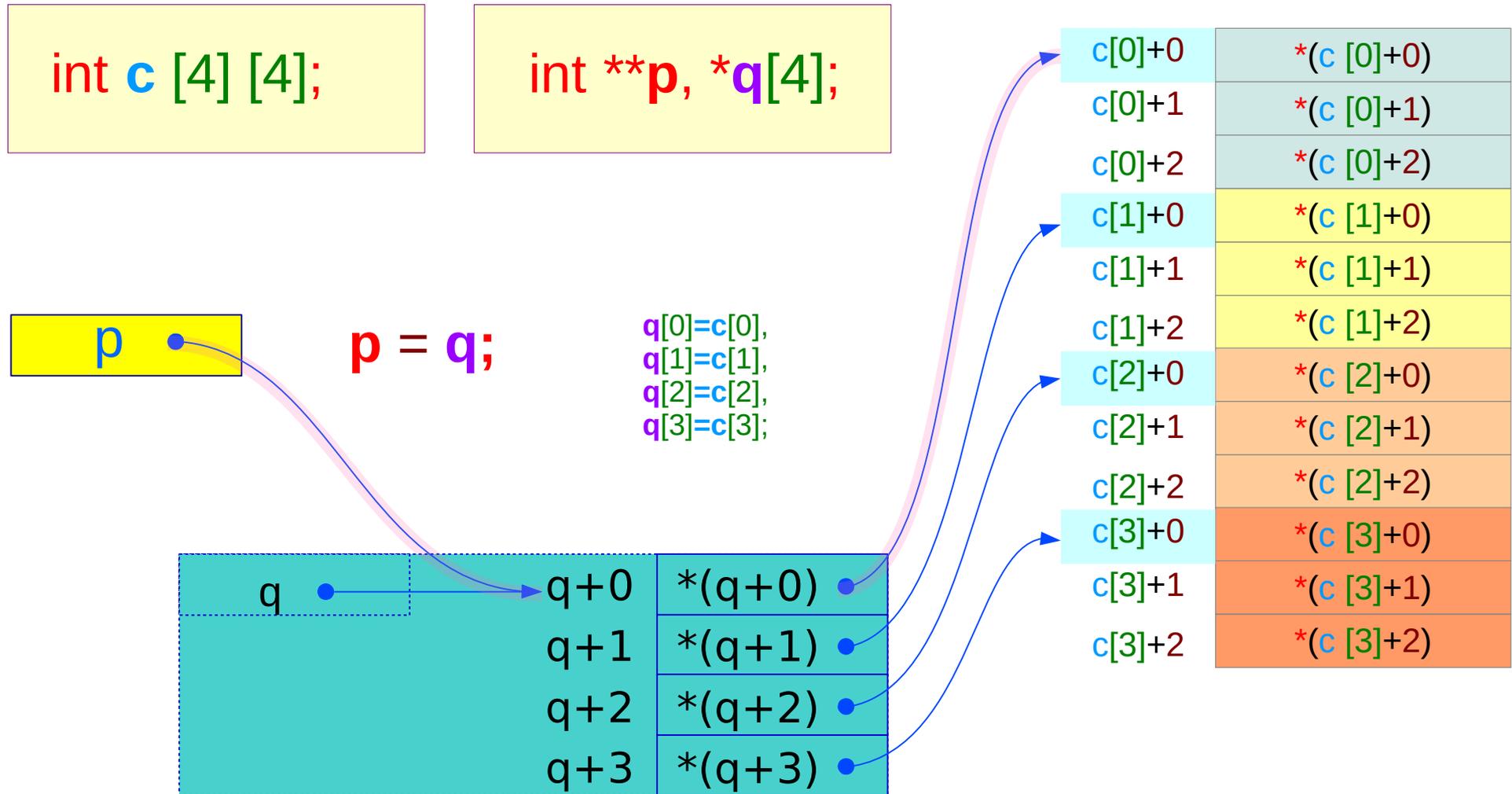


$(*(p+i))[j]$



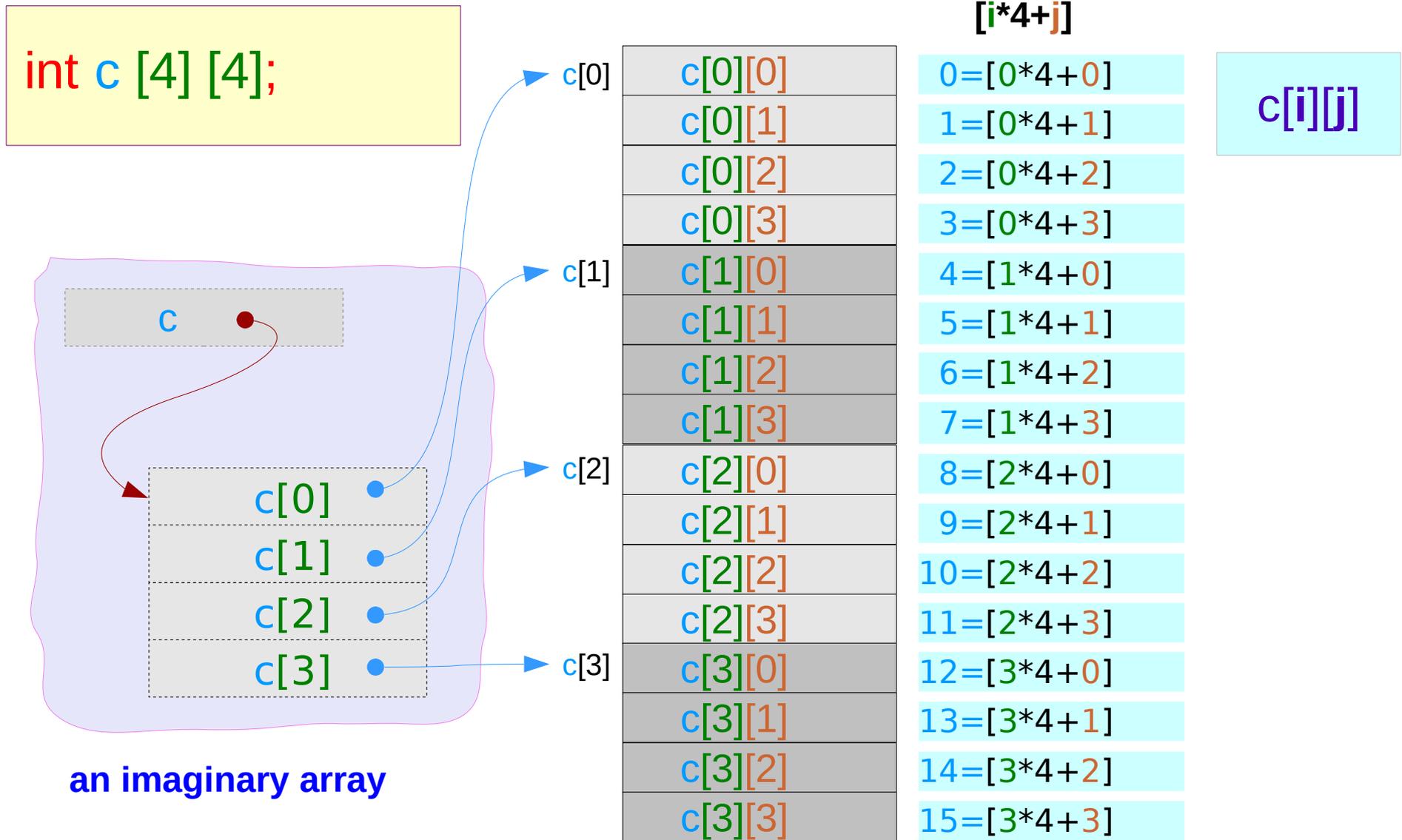
$*(*(p+i)+j)$

2-d array access via a double pointer



to use a double pointer `p`, a pointer array `q` must be allocated and initialized

A 2-d array stored as a 1-d array (row major order)



A Linear Memory Address

1-d address

c[0]

$$c[0] + 0 = c[0] + 0 \cdot 4 + 0$$



$$c[0] + 1 = c[0] + 0 \cdot 4 + 1$$



$$c[0] + 2 = c[0] + 0 \cdot 4 + 2$$



$$c[0] + 3 = c[0] + 0 \cdot 4 + 3$$



$$c[0] + 4 = c[0] + 1 \cdot 4 + 0$$



$$c[0] + 5 = c[0] + 1 \cdot 4 + 1$$



$$c[0] + 6 = c[0] + 1 \cdot 4 + 2$$



$$c[0] + 7 = c[0] + 1 \cdot 4 + 3$$



$$c[0] + 8 = c[0] + 2 \cdot 4 + 0$$



$$c[0] + 9 = c[0] + 2 \cdot 4 + 1$$



$$c[0] + 10 = c[0] + 2 \cdot 4 + 2$$



$$c[0] + 11 = c[0] + 2 \cdot 4 + 3$$



$$c[0] + 12 = c[0] + 3 \cdot 4 + 0$$



$$c[0] + 13 = c[0] + 3 \cdot 4 + 1$$



$$c[0] + 14 = c[0] + 3 \cdot 4 + 2$$



$$c[0] + 15 = c[0] + 3 \cdot 4 + 3$$



2-d address

c[0],c[1],c[2],c[3]

$$c[0] + 0 = *(c + 0) + 0$$

(int) c[0][0]

$$c[0] + 1 = *(c + 0) + 1$$

(int) c[0][1]

$$c[0] + 2 = *(c + 0) + 2$$

(int) c[0][2]

$$c[0] + 3 = *(c + 0) + 3$$

(int) c[0][3]

$$c[1] + 0 = *(c + 1) + 0$$

(int) c[1][0]

$$c[1] + 1 = *(c + 1) + 1$$

(int) c[1][1]

$$c[1] + 2 = *(c + 1) + 2$$

(int) c[1][2]

$$c[1] + 3 = *(c + 1) + 3$$

(int) c[1][3]

$$c[2] + 0 = *(c + 2) + 0$$

(int) c[2][0]

$$c[2] + 1 = *(c + 2) + 1$$

(int) c[2][1]

$$c[2] + 2 = *(c + 2) + 2$$

(int) c[2][2]

$$c[2] + 3 = *(c + 2) + 3$$

(int) c[2][3]

$$c[3] + 0 = *(c + 3) + 0$$

(int) c[3][0]

$$c[3] + 1 = *(c + 3) + 1$$

(int) c[3][1]

$$c[3] + 2 = *(c + 3) + 2$$

(int) c[3][2]

$$c[3] + 3 = *(c + 3) + 3$$

(int) c[3][3]

A linearization of a 2-D array

1-d view

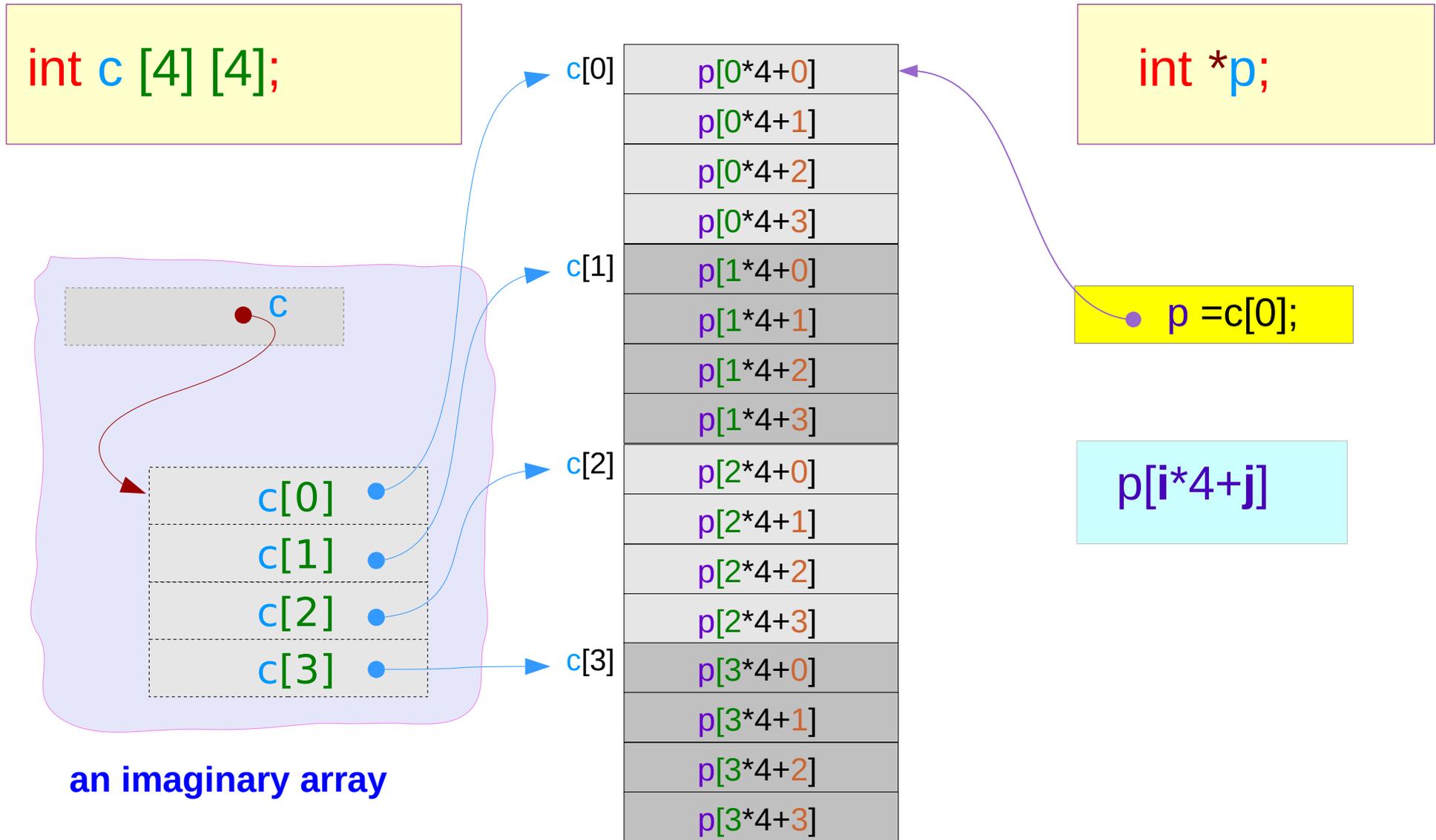
p+0	(int) p[0]
p+1	(int) p[1]
p+2	(int) p[2]
p+3	(int) p[3]
p+4	(int) p[4]
p+5	(int) p[5]
p+6	(int) p[6]
p+7	(int) p[7]
p+8	(int) p[8]
p+9	(int) p[9]
p+10	(int) p[10]
p+11	(int) p[11]
p+12	(int) p[12]
p+13	(int) p[13]
p+14	(int) p[14]
p+15	(int) p[15]



2-d view

$c[0]+0 = *(c+0)+0$	(int) c[0][0]
$c[0]+1 = *(c+0)+1$	(int) c[0][1]
$c[0]+2 = *(c+0)+2$	(int) c[0][2]
$c[0]+3 = *(c+0)+3$	(int) c[0][3]
$c[1]+0 = *(c+1)+0$	(int) c[1][0]
$c[1]+1 = *(c+1)+1$	(int) c[1][1]
$c[1]+2 = *(c+1)+2$	(int) c[1][2]
$c[1]+3 = *(c+1)+3$	(int) c[1][3]
$c[2]+0 = *(c+2)+0$	(int) c[2][0]
$c[2]+1 = *(c+2)+1$	(int) c[2][1]
$c[2]+2 = *(c+2)+2$	(int) c[2][2]
$c[2]+3 = *(c+2)+3$	(int) c[2][3]
$c[3]+0 = *(c+3)+0$	(int) c[3][0]
$c[3]+1 = *(c+3)+1$	(int) c[3][1]
$c[3]+2 = *(c+3)+2$	(int) c[3][2]
$c[3]+3 = *(c+3)+3$	(int) c[3][3]

2-d array access via a single pointer



2-d array index vs 1-d array index

```
int c [4] [4];
```

```
int *p=c[0];
```

$c[i][j]$

$p[i*4+j]$

c[0]	c[0][0]
	c[0][1]
	c[0][2]
	c[0][3]
c[1]	c[1][0]
	c[1][1]
	c[1][2]
	c[1][3]
c[2]	c[2][0]
	c[2][1]
	c[2][2]
	c[2][3]
c[3]	c[3][0]
	c[3][1]
	c[3][2]
	c[3][3]

p[0*4+0]
p[0*4+1]
p[0*4+2]
p[0*4+3]
p[1*4+0]
p[1*4+1]
p[1*4+2]
p[1*4+3]
p[2*4+0]
p[2*4+1]
p[2*4+2]
p[2*4+3]
p[3*4+0]
p[3*4+1]
p[3*4+2]
p[3*4+3]

2-d array access via a single pointer

```
int *p = c[0];
```



```
int c [4][4];
```

```
p[ i*4 + j ]
```



```
c[ i ][ j ]
```

```
*(p + i*4 + j)
```



```
*(*(c+i) + j)
```

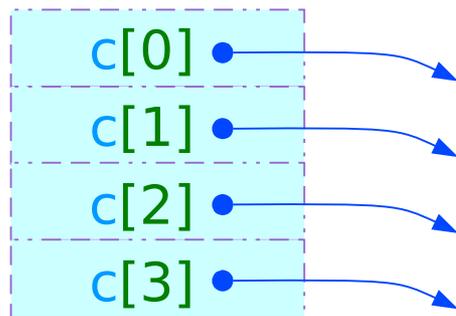
```
*(p + k)    i = k / 4;  
            j = k % 4;
```

Static Allocation of a 2-d Array

```
int A [4][3];
```

A in %eax,
i in %edx,
j in %ecx

```
sall    $2, %ecx           ;; j * 4  
leal   (%edx, %edx, 2), %edx  ;; i * 3  
leal   (%ecx, %edx, 4), %edx  ;; j * 4 + i * 12  
movl   (%eax, %edx), %eax     ;; read M[ XA+4(3i +j) ]
```



The pointer array :
not allocated
in the memory

c[0]+0	*(c [0]+0)
c[0]+1	*(c [0]+1)
c[0]+2	*(c [0]+2)
c[1]+0	*(c [1]+0)
c[1]+1	*(c [1]+1)
c[1]+2	*(c [1]+2)
c[2]+0	*(c [2]+0)
c[2]+1	*(c [2]+1)
c[2]+2	*(c [2]+2)
c[3]+0	*(c [3]+0)
c[3]+1	*(c [3]+1)
c[3]+2	*(c [3]+2)

The name of a 2-d array

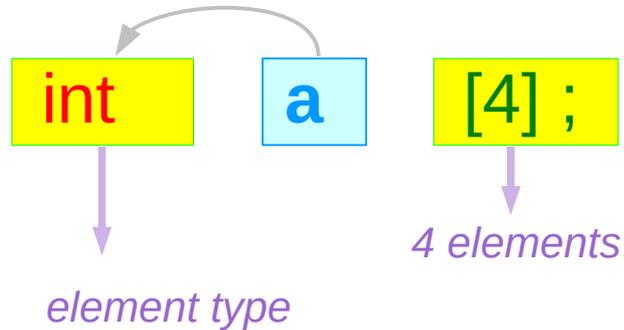
```
int    a [4];
```

```
int    c [4] [4];
```

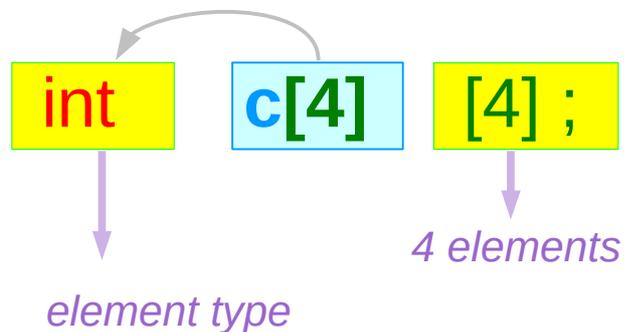
1. the name of the nested array (recursive definition)
2. a double pointer
3. a pointer to an array

c: the nested array name

```
int    a [4];  
  
int    c [4] [4];
```



The array name **a** holds the starting address of the 4 integer element array



c[0], c[1], c[2], c[3] holds the starting address of the 4 integer element array

The array name **c** is the name of the nested array

c is a double pointer and a **1-d** array pointer

```
int    a [4];
```

```
int    c [4] [4];
```

*****(*****(**c**+**0**)+**0**)



****c**

a double pointer

(*****(**c**+**0**))[**j**]



(***c**)[**j**]

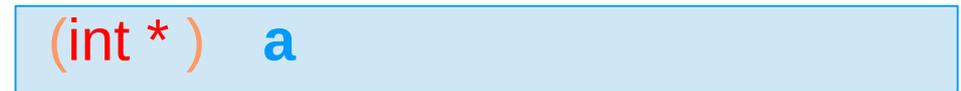
a **1-d** array pointer

c is a double pointer

```
int    a [4];  
int    c [4] [4];
```

****c**

***(* (c+0)+0)**



a points to an integer data



c[i] points to an integer data
c points to an integer pointer

c is also a pointer to an array

```
int    a [4];  
int    c [4] [4];
```

(*c) [j]
(*(c+0)) [j]



a : an array name



c[i] : array names
c : the pointer array name

Pointer to Arrays

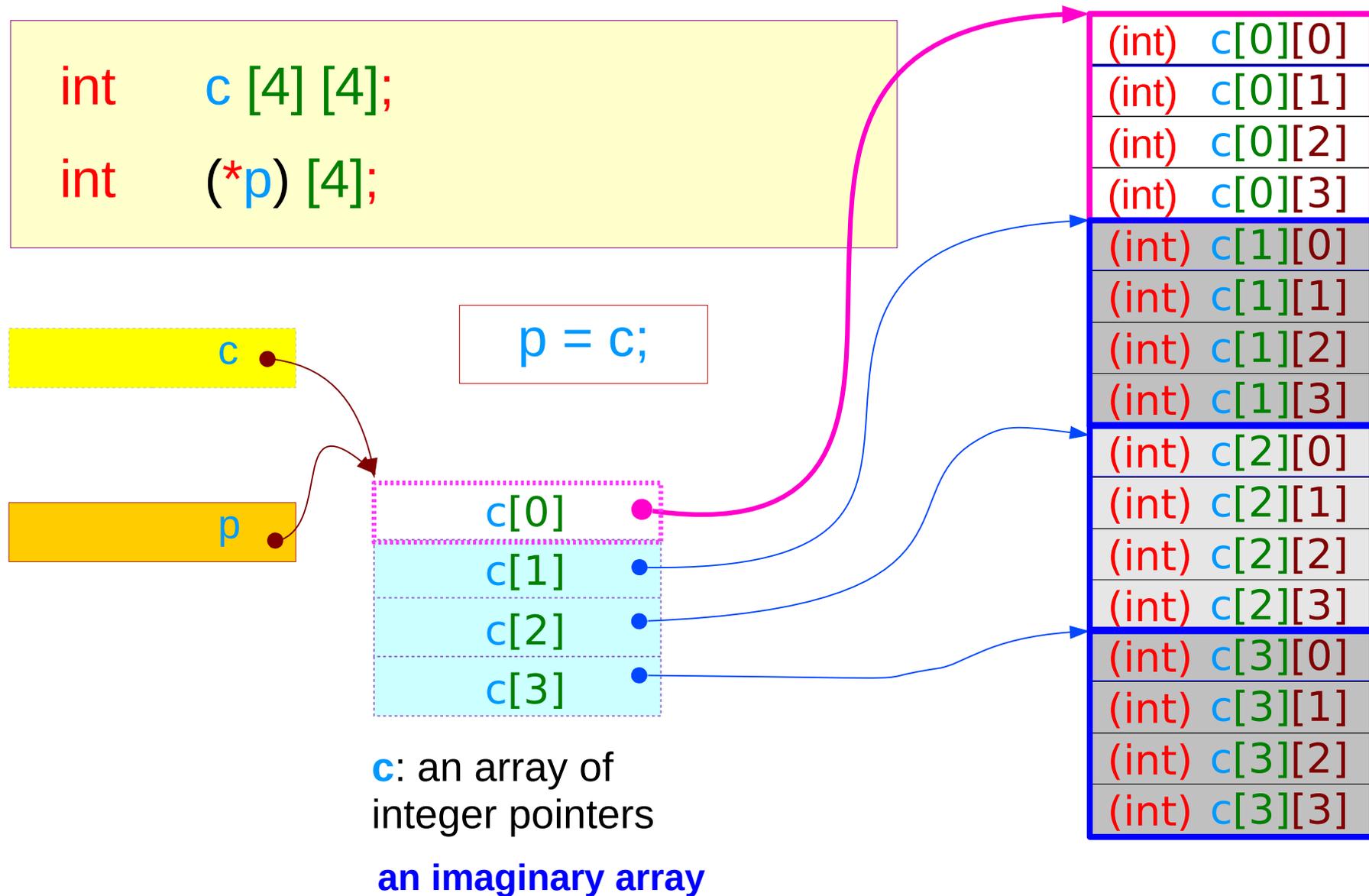
1. pointer to 1-d arrays

```
int (*p) [4];
```

2. pointer to 2-d arrays

```
int (*p) [4][4];
```

Pointer to 1-d arrays



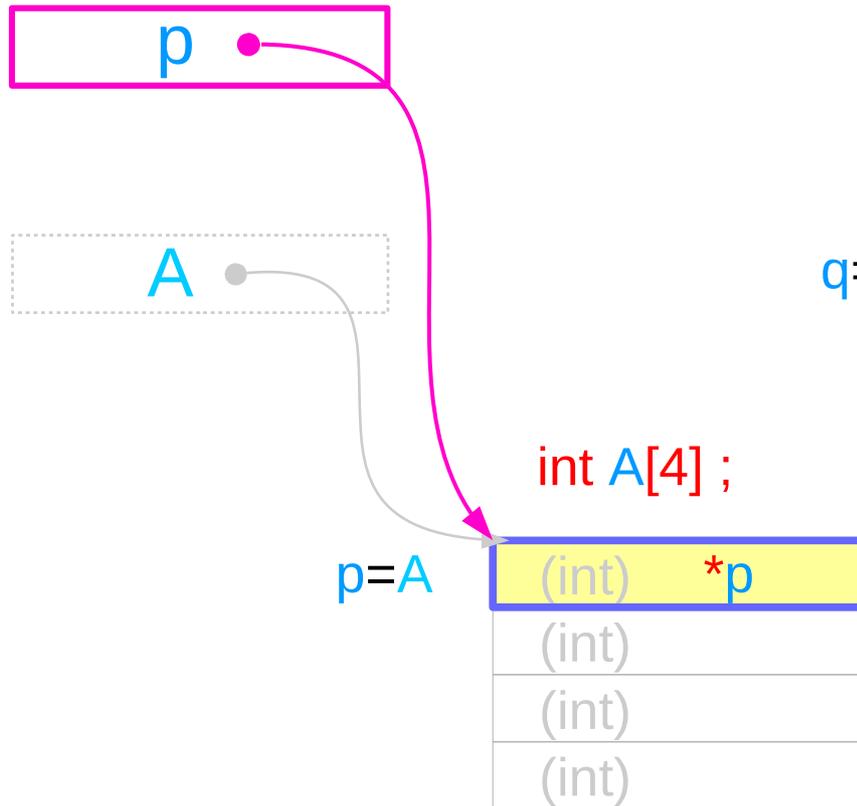
Pointer to Integer vs. Pointer to Array

```
int *p ;
```

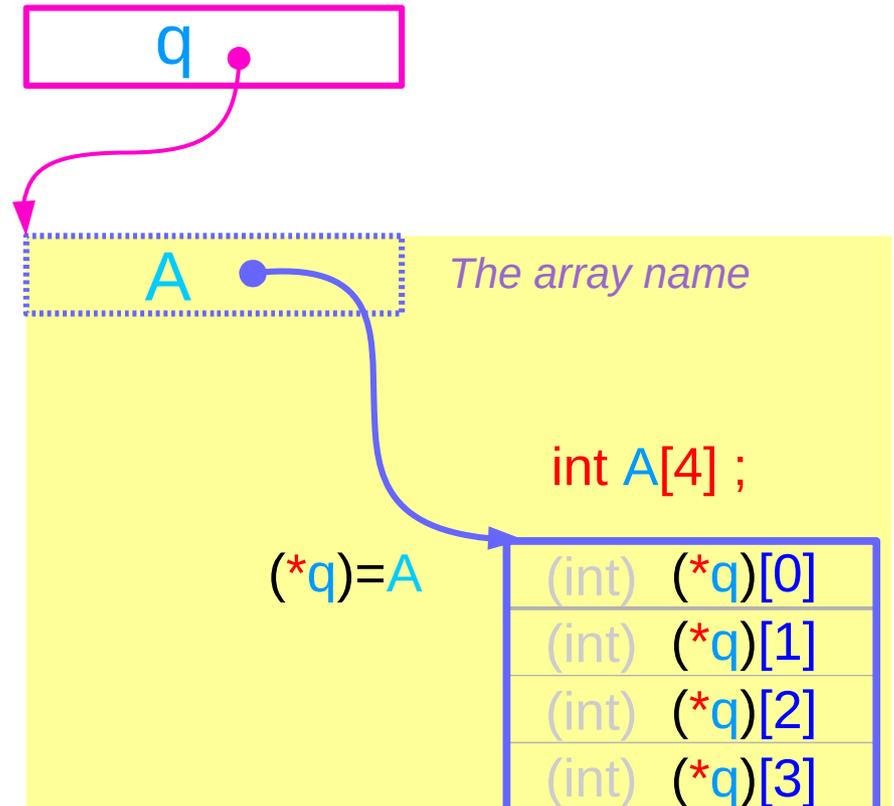
```
p = A ;
```

```
int (*q) [4];
```

```
q = &A ;
```



The int pointer type



The array type

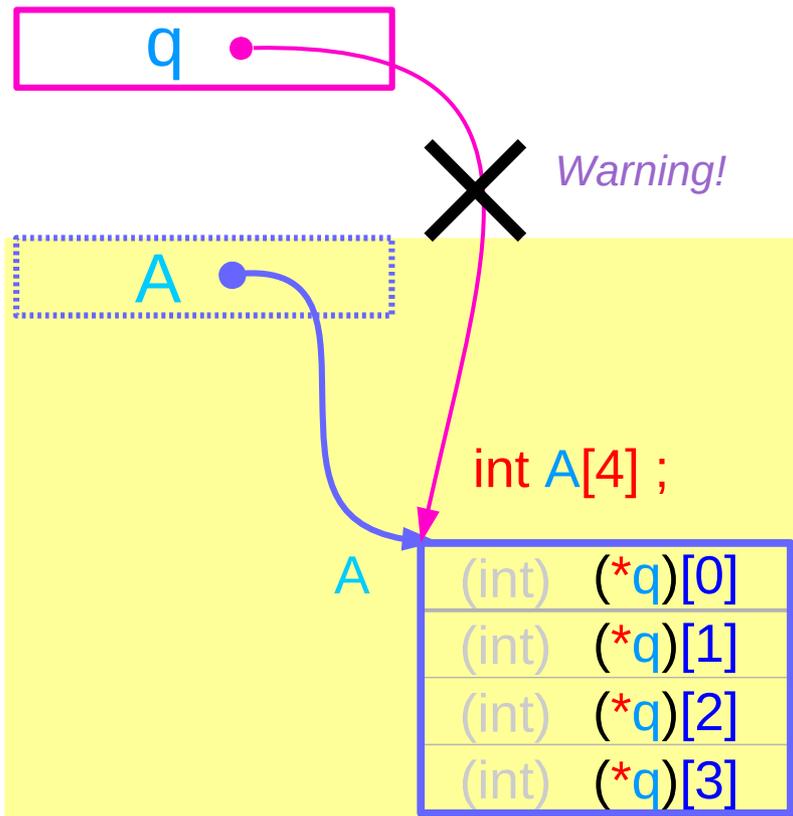
Must point to an array type (array name)

```
int (*q) [4];
```

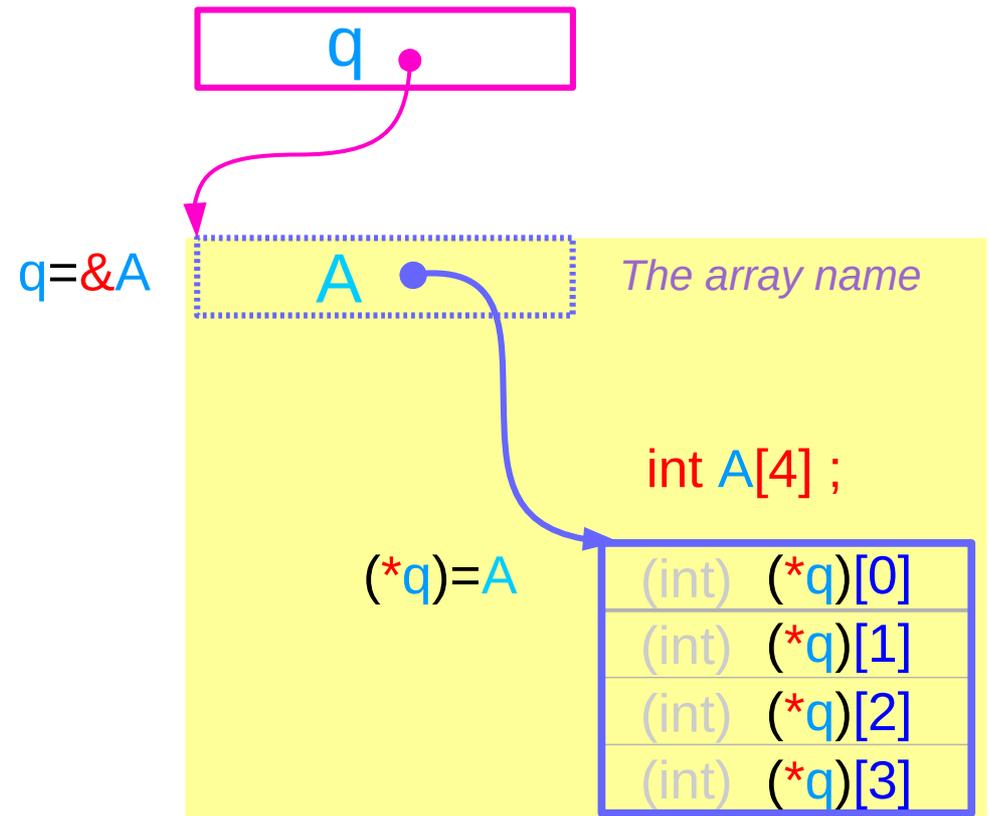
```
q = A ;
```

```
int (*q) [4];
```

```
q = &A ;
```



The array type



The array type

Pointer to a 2-d array

```
int c [4][4];  
int (*p) [4][4] = &c;  
int (*q) [4]    = &c[0] ;
```

2-d array c

2-d array pointer p

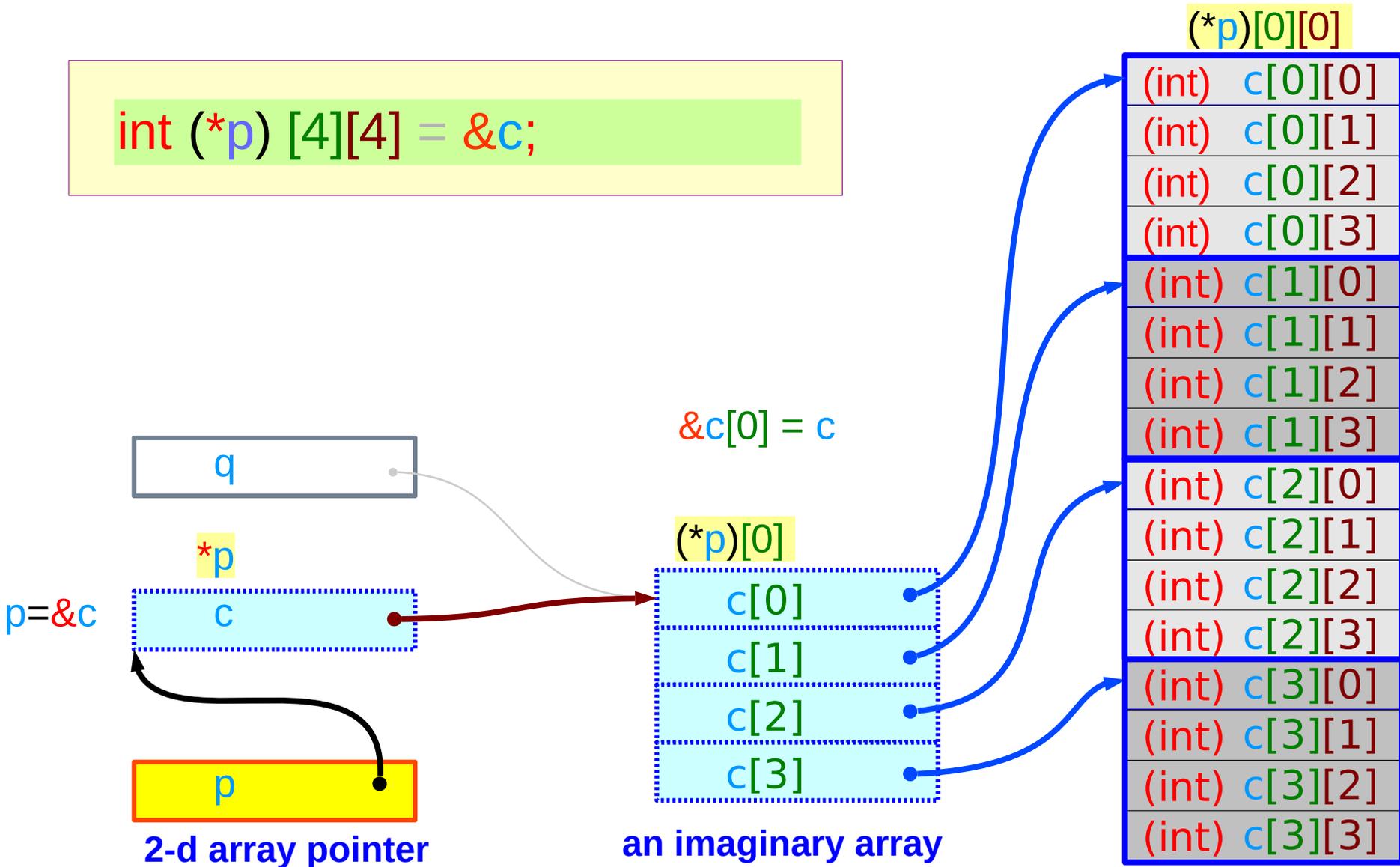
1-d array pointer q

(*p)[i][j] → c[i][j]

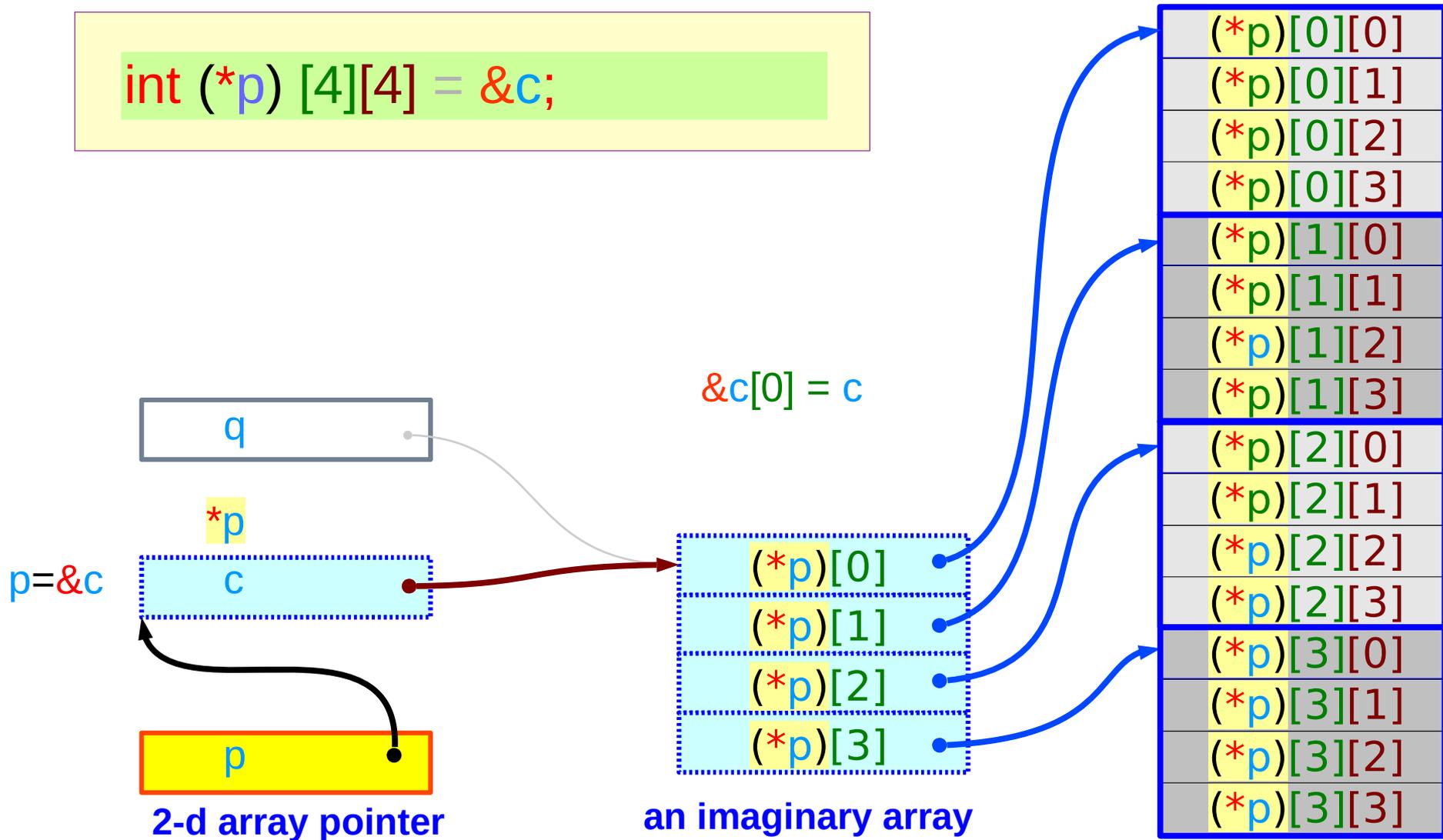
q[i][j] → c[i][j]

Pointer to a 2-d array

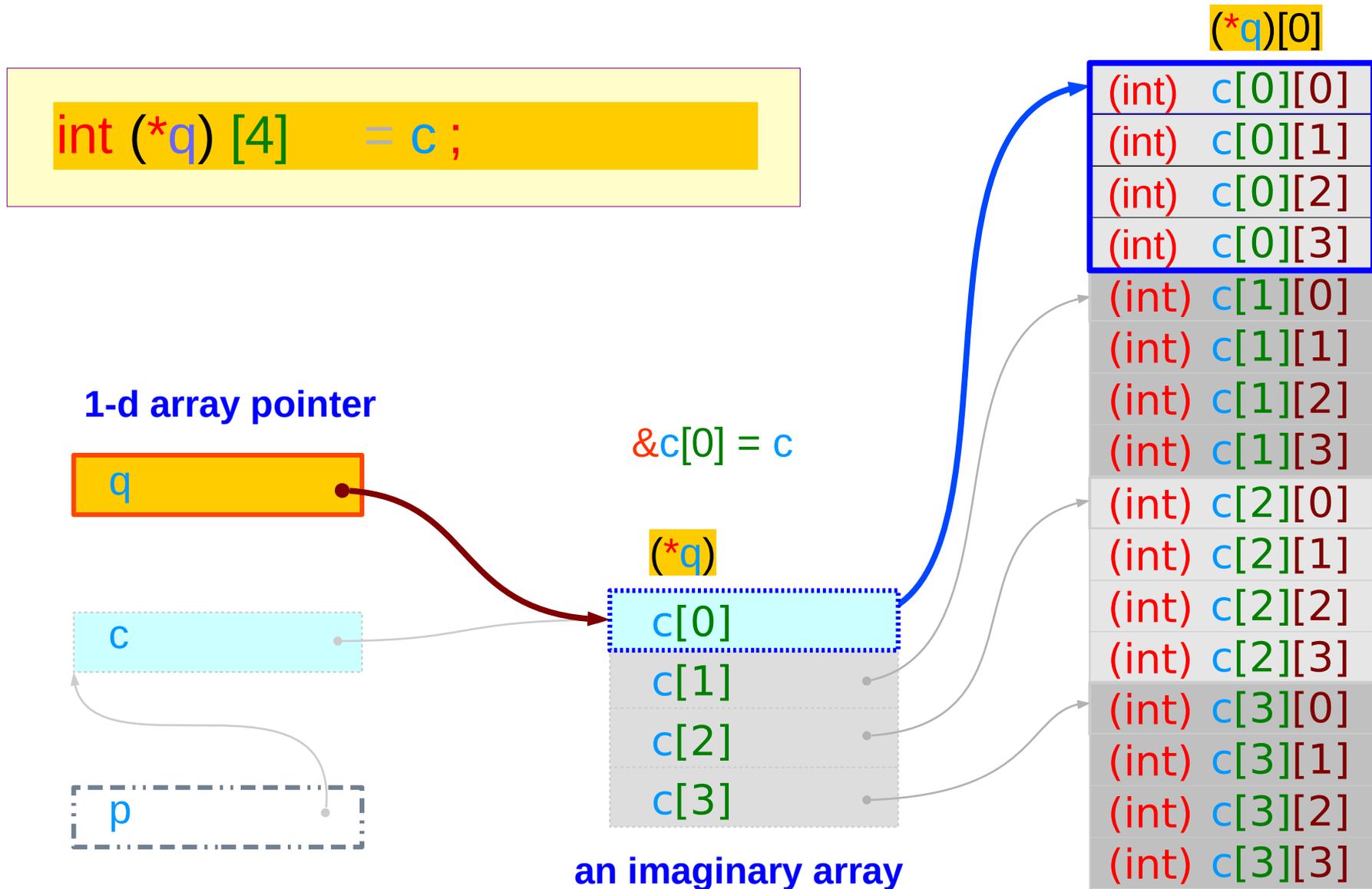
```
int (*p) [4][4] = &c;
```



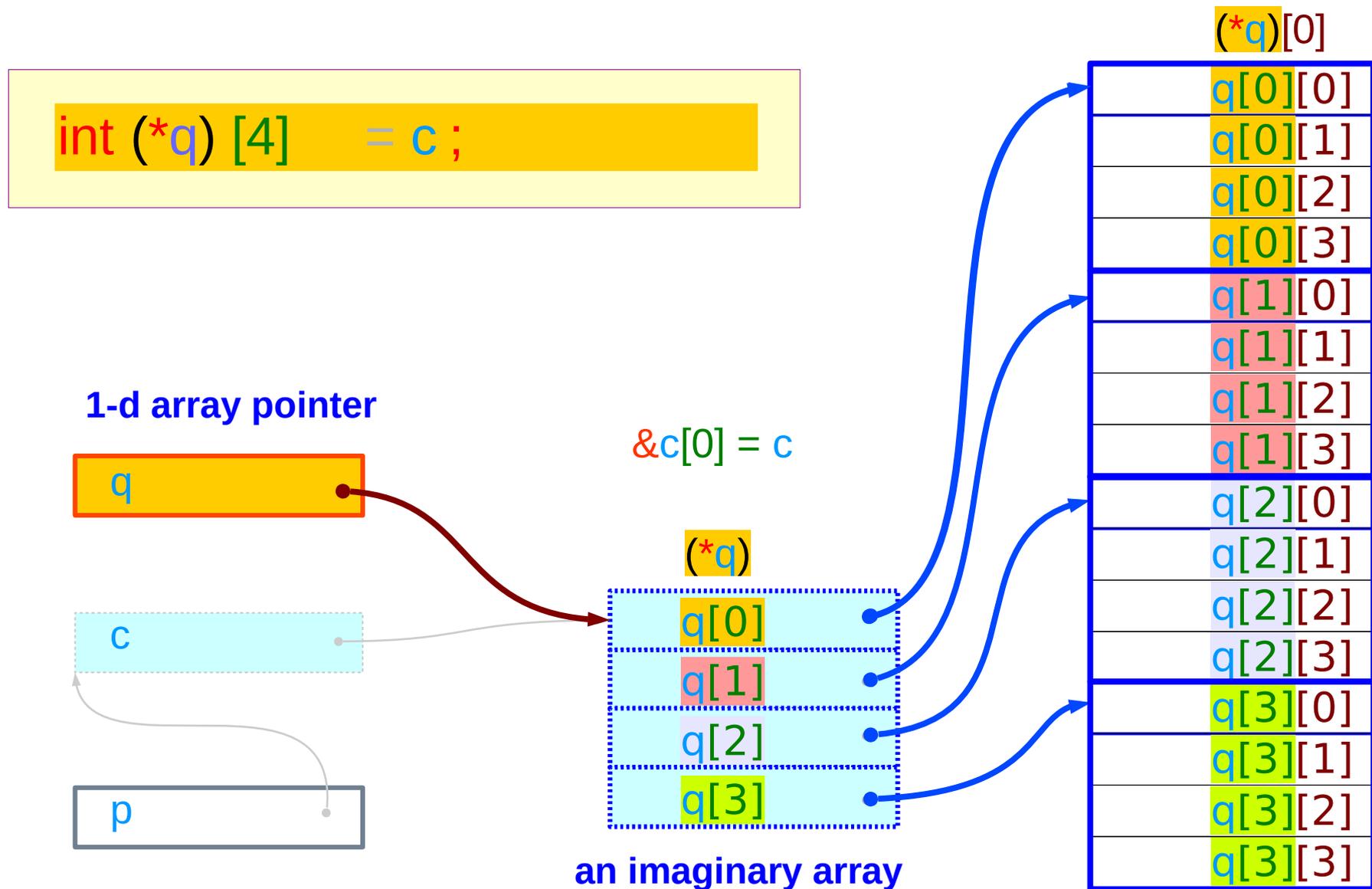
2-d array access using a pointer to a 2-d array



Pointer to a 1-d array

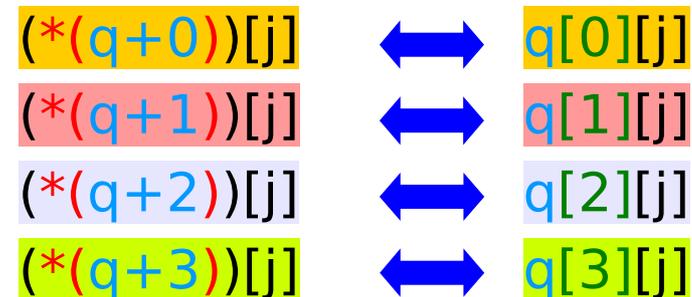
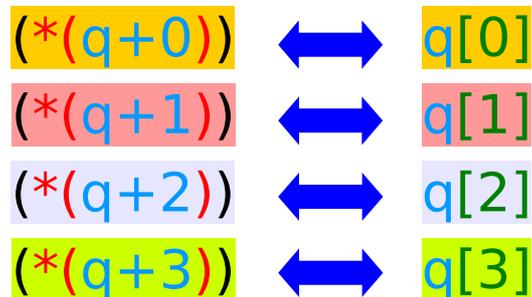
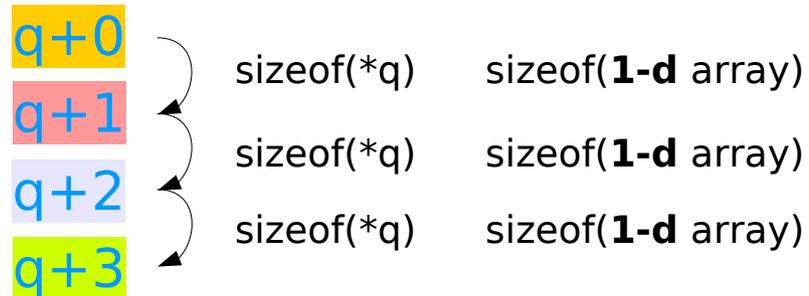


2-d array access using a pointer to a 1-d array



Incrementing a pointer to a **1-d** array

```
int (*q) [4] = c;
```



Dynamic Memory Allocation of 2-d Arrays

1. method 1

```
int ** c ;  
c = malloc(4 * sizeof (int *) ) ;  
c[0] = malloc(4 * sizeof (int) ) ;  
c[1] = malloc(4 * sizeof (int) ) ;  
c[2] = malloc(4 * sizeof (int) ) ;  
c[3] = malloc(4 * sizeof (int) ) ;
```

2. method 2

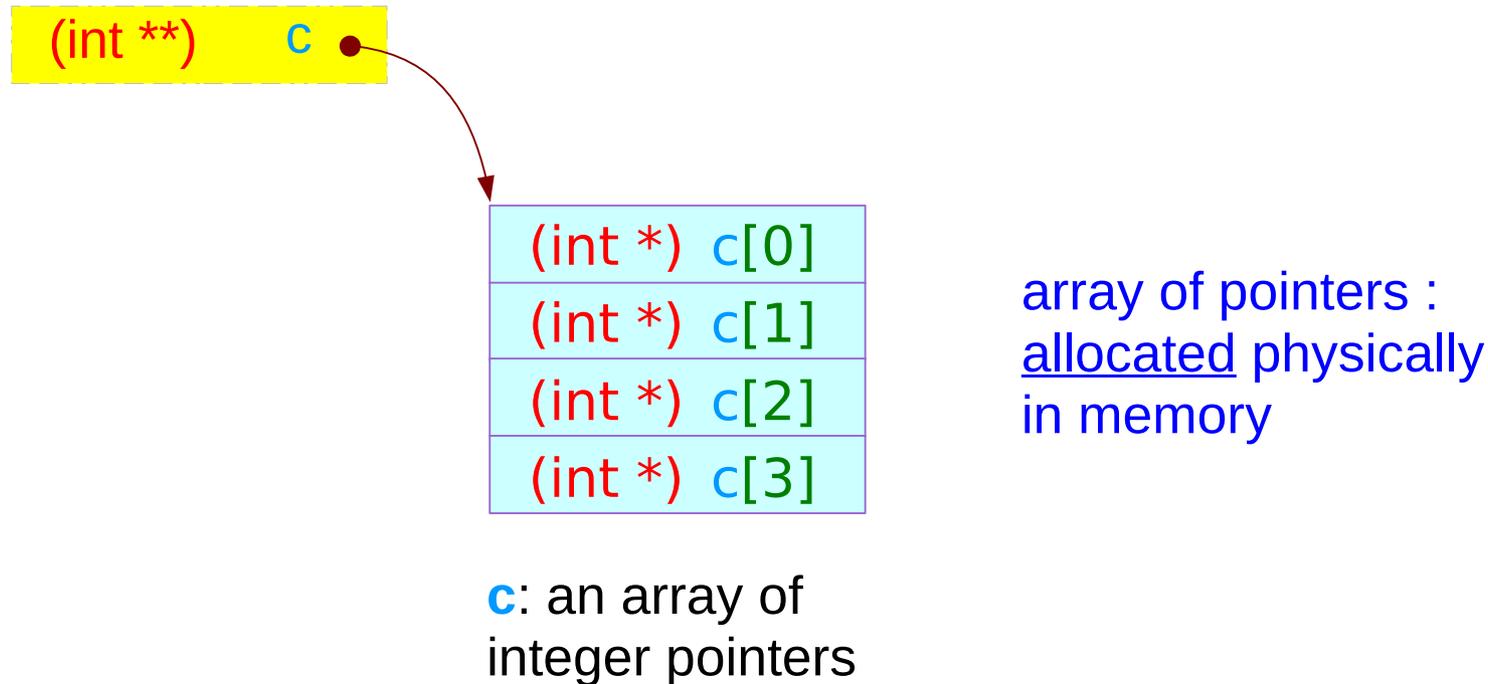
```
int (*p) [4] ;  
p = malloc(4 * 4 * sizeof (int)) ;
```

3. method 3

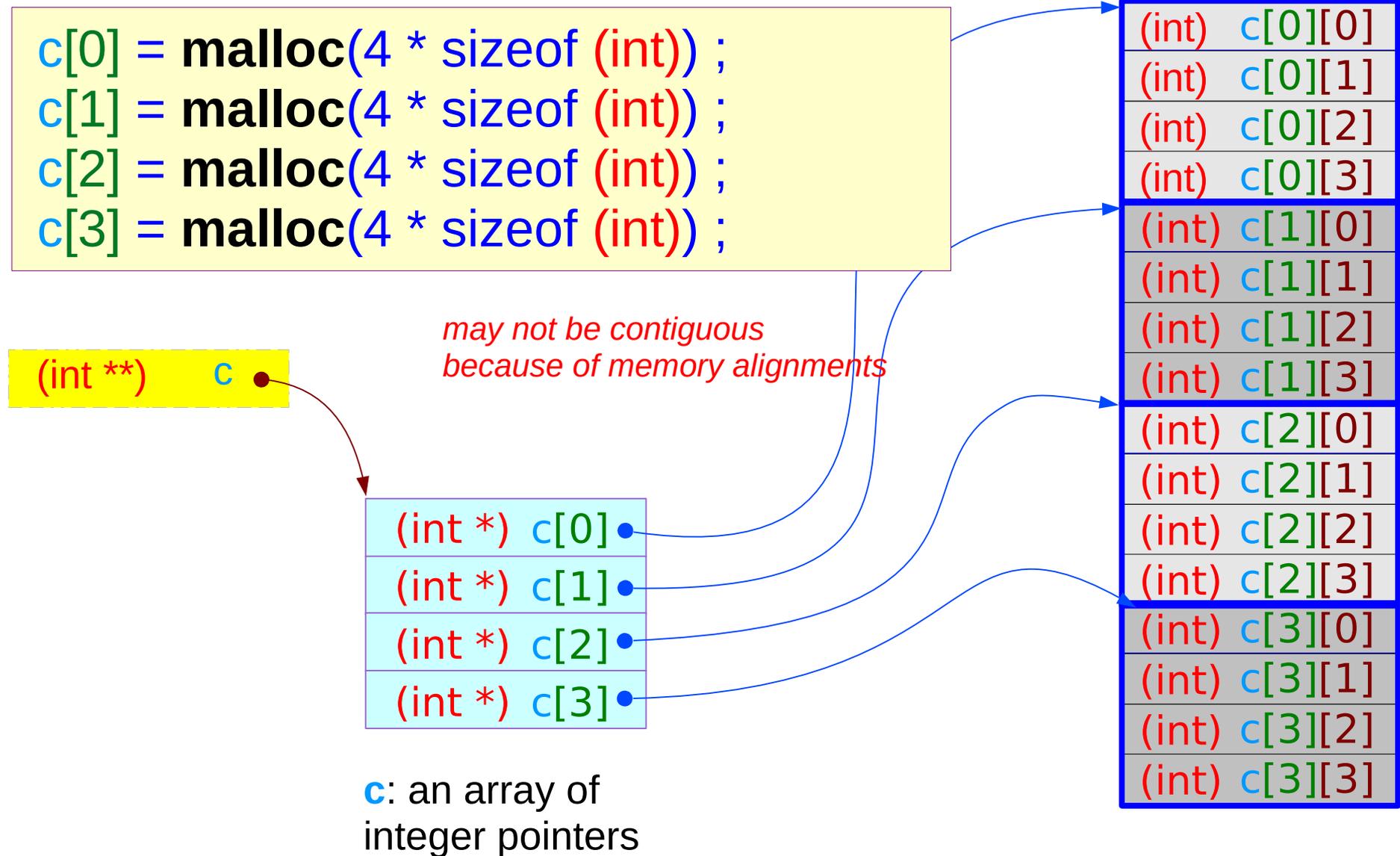
```
int ** c ;  
int * p ;  
c = malloc( 4 * sizeof(int *) ) ;  
p = malloc( 4 * 4 * sizeof(int) ) ;  
for (i=0; i<M; i++) c[i] = p + i*N;
```

2-d array dynamic allocation : method 1 (a)

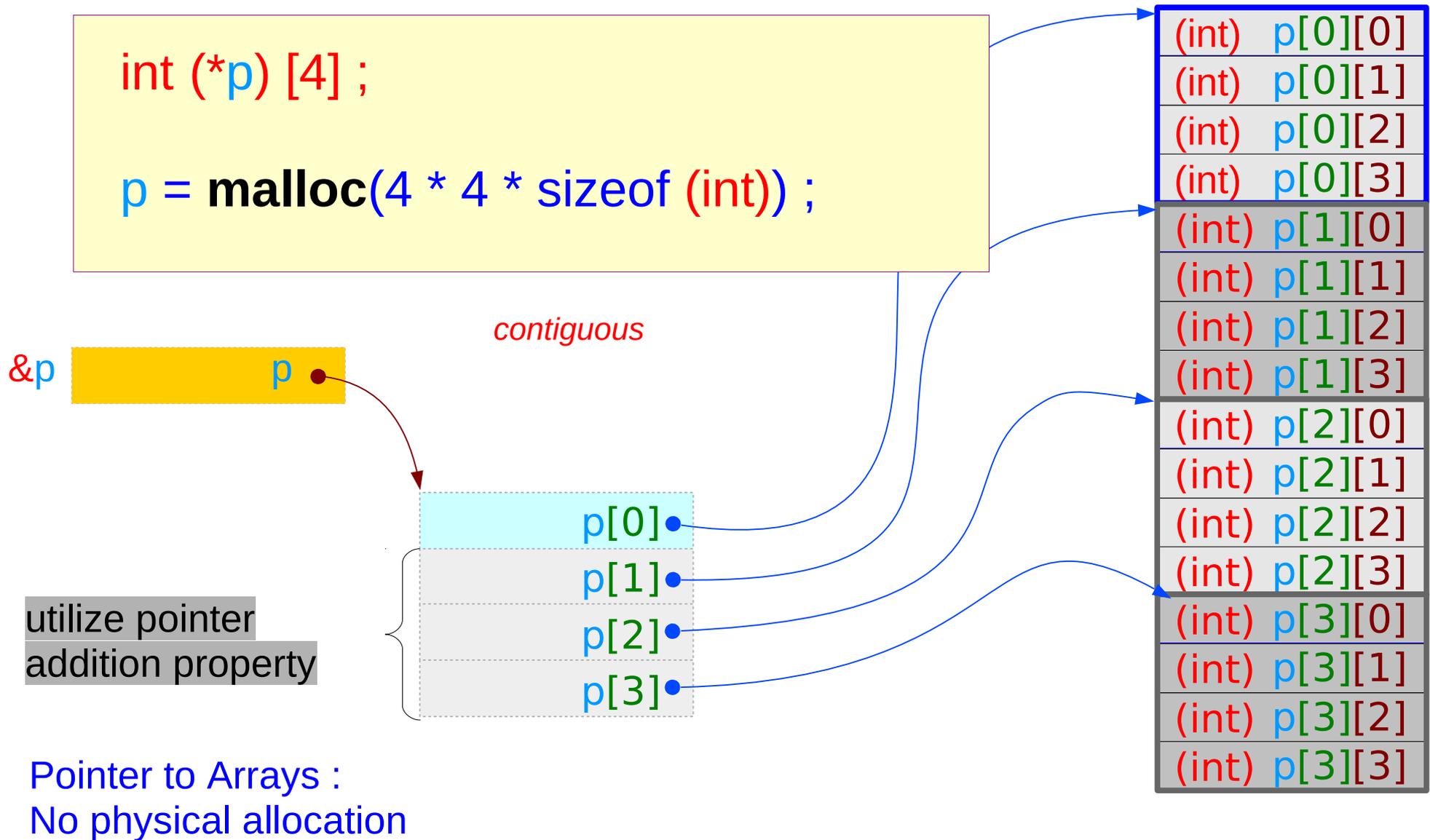
```
int ** c ;  
  
c = malloc(4 * sizeof (int *)) ;
```



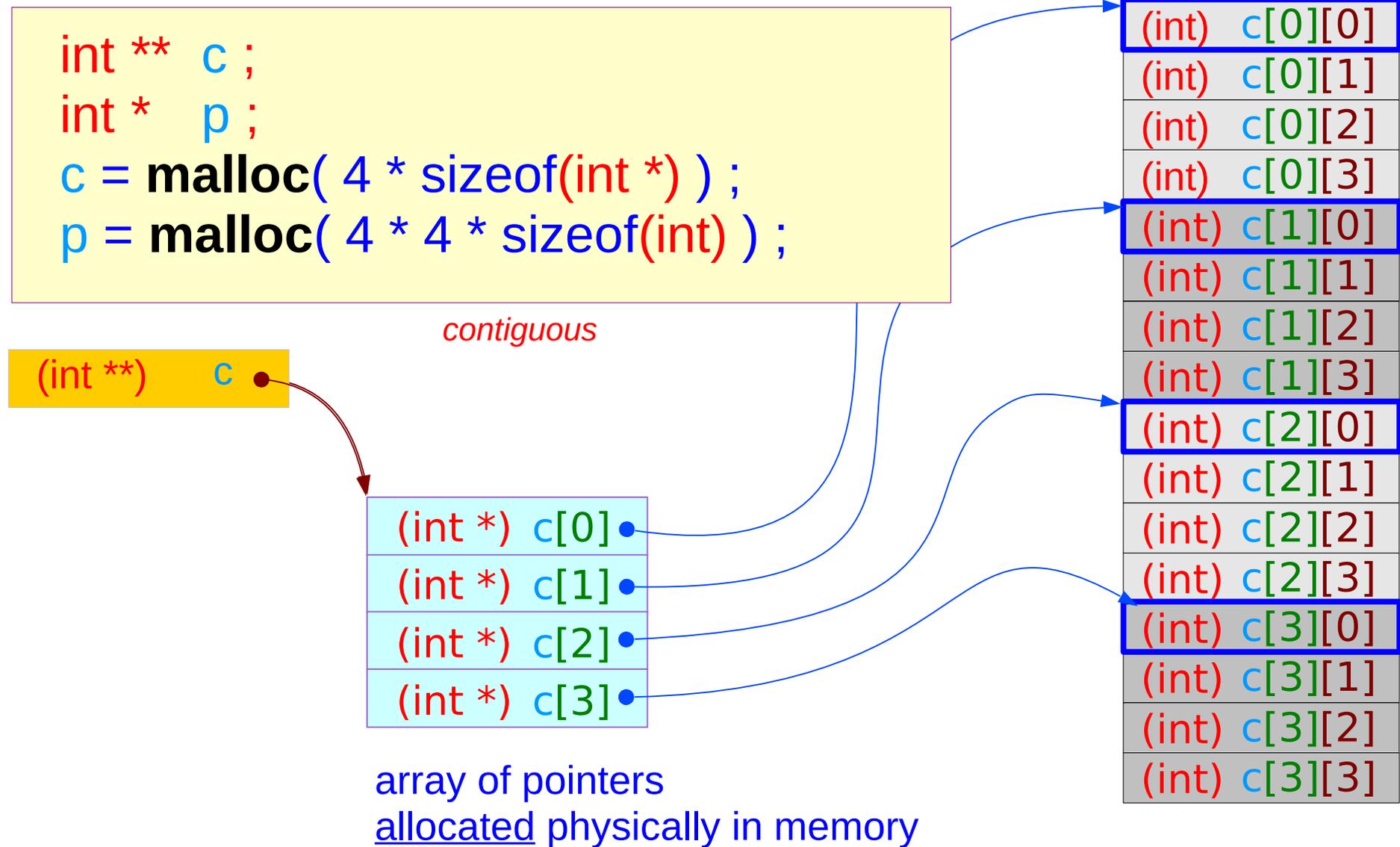
2-d array dynamic allocation : method 1 (b)



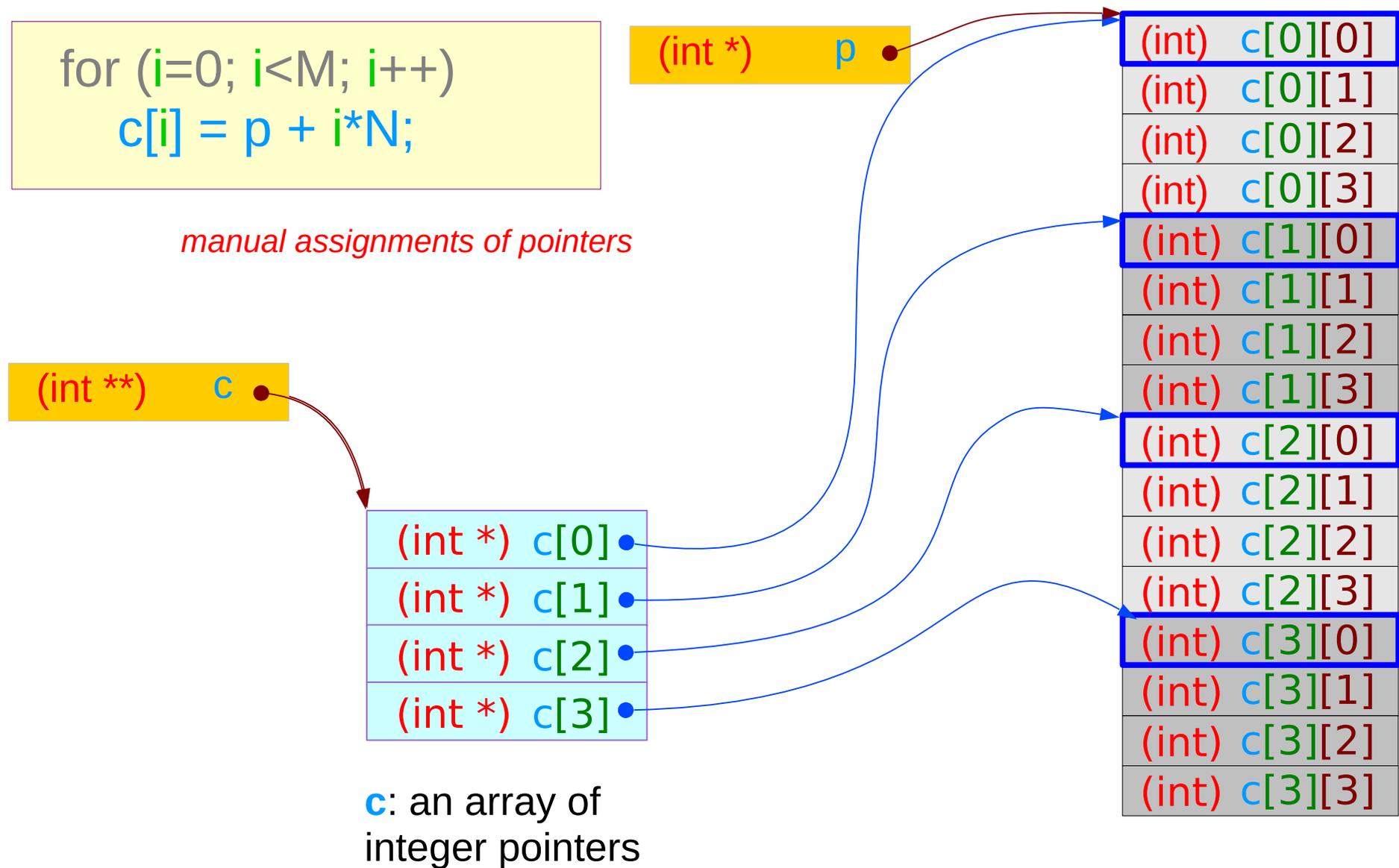
2-d array dynamic allocation : method 2



2-d array dynamic allocation : method 3 (a)



2-d array dynamic allocation : method 3 (b)



Limitations

No index Range Checking

Array Size must be a constant expression

Variable Array Size

Arrays cannot be Copied or Compared

Aggregate Initialization and Global Arrays

Precedence Rule

Index Type Must be Integral

3-d Array Index

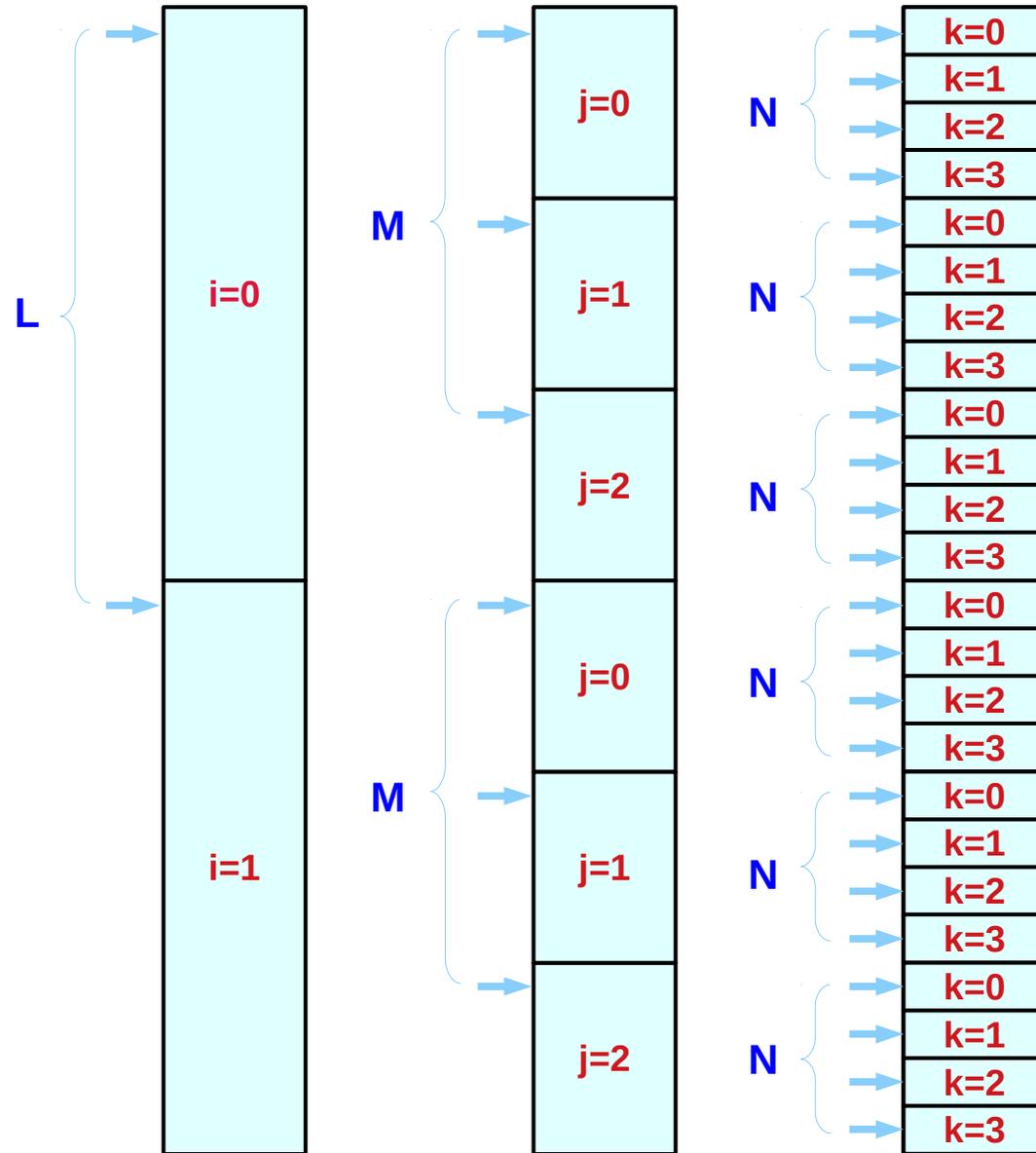
c [i][j][k]

interpret as recursive indirections
interpret as hierarchical sub-arrays

L, M, N – the number of index values

```
int c [L][M][N];
```

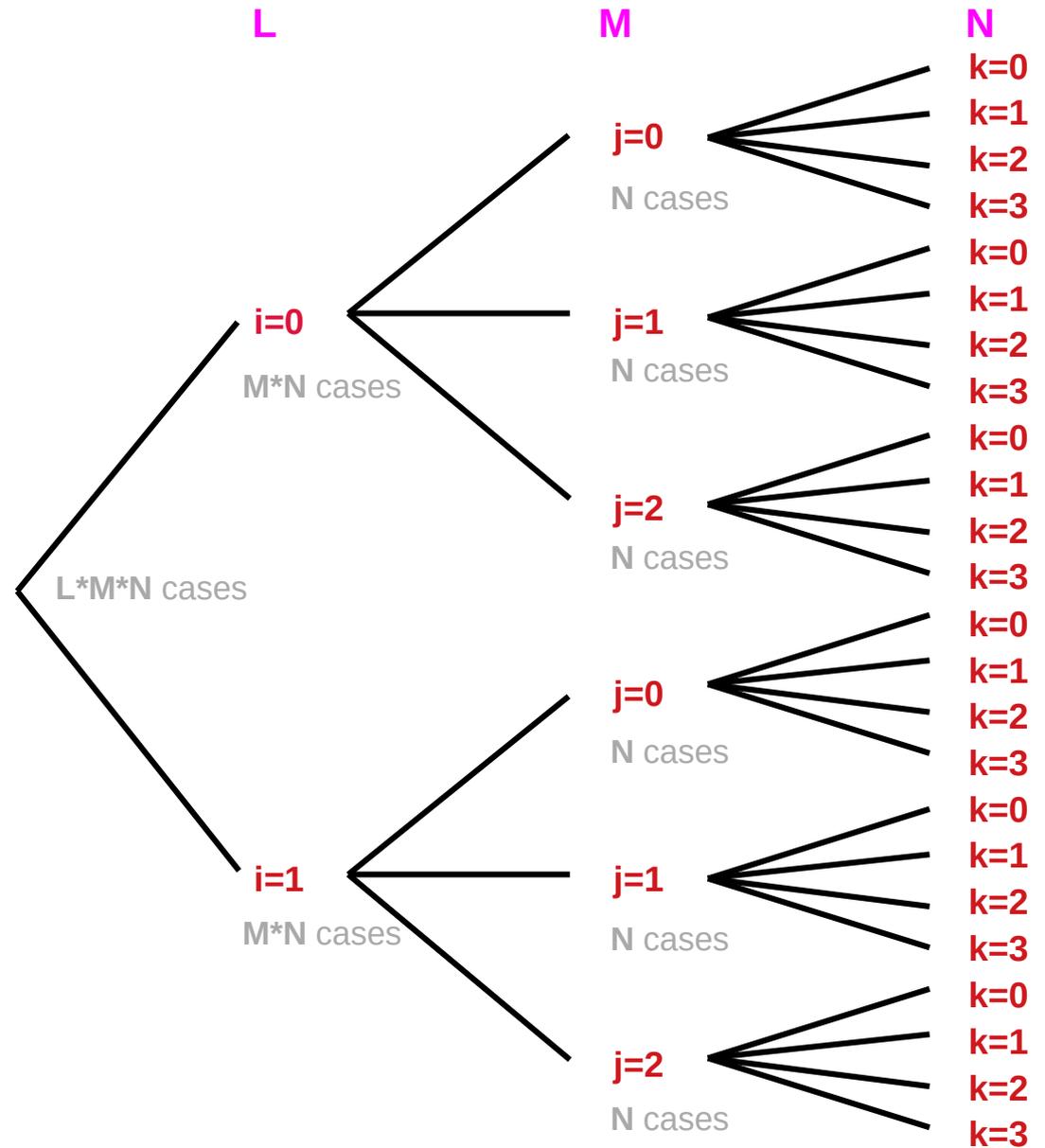
```
c [i][j][k]
```



Index value tree – all possible combinations

```
int c [L][M][N];
```

```
c [i][j][k]
```



The number of elements of subarrays

```
int c [L][M][N];
```

```
c [i][j][k]
```

c **L*M*N** elements

c[i] **M*N** elements

c[i][j] **N** elements



array
names

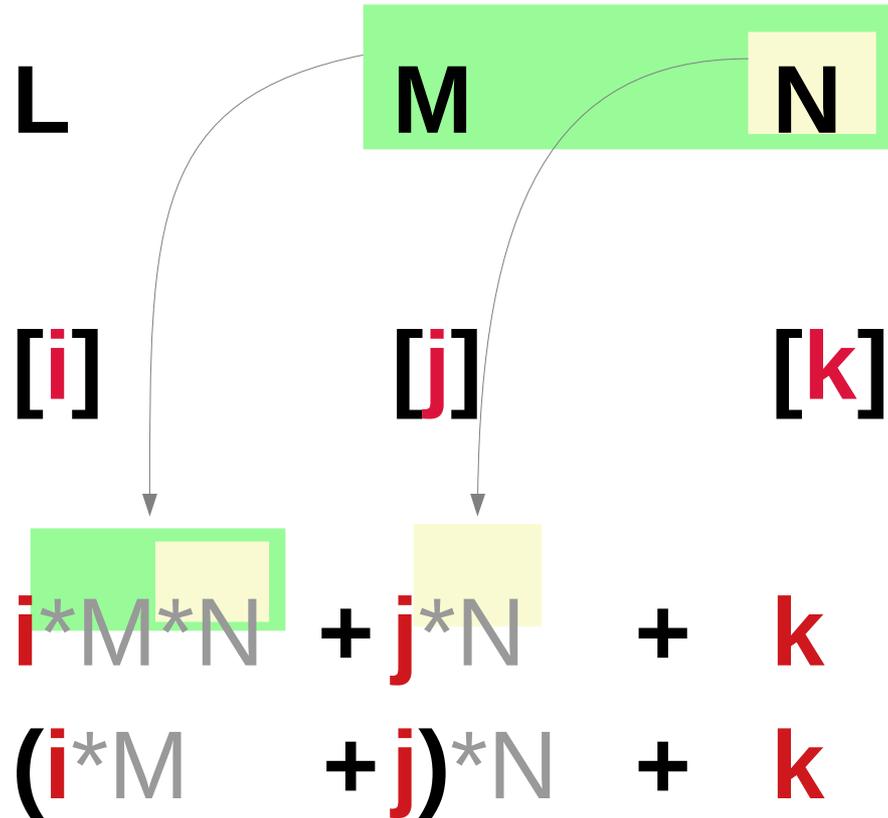


covering
elements

From a 3-d index to a 1-d index

```
int c [L][M][N];
```

```
c [i][j][k]
```



$i * M * N, j * N, k$ – index offset values

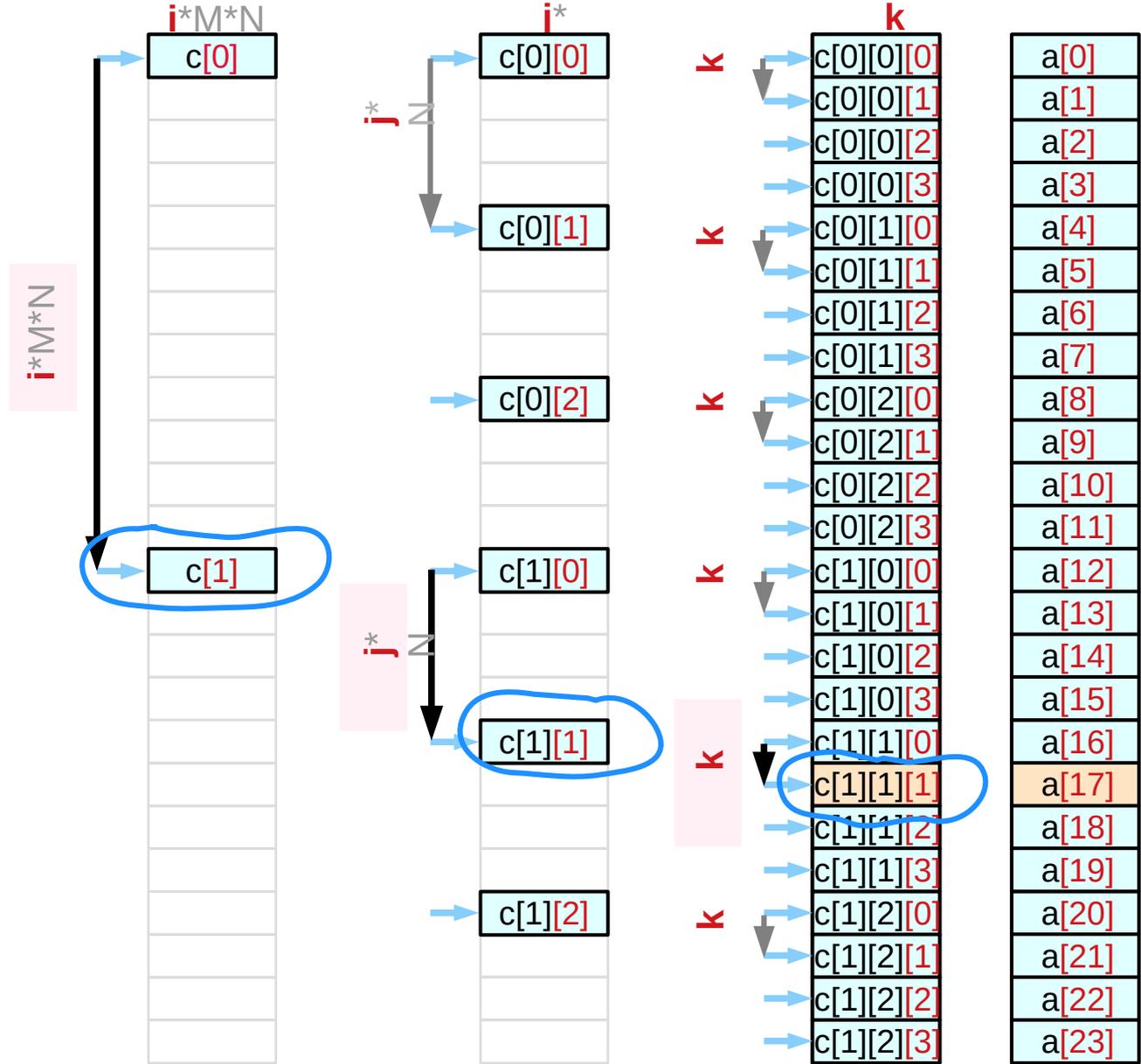
```
int c[L][M][N];
```

```
c[i][j][k]
```

```
c[1][1][1]
```

$i=1$	$j=1$	$k=1$
-------	-------	-------

$$a[(1 * 3 + 1) * 4 + 1]$$



References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun
- [5] <https://pdos.csail.mit.edu/6.828/2008/readings/pointers.pdf>